Spatiotemporal Computational Mechanics

Paul Riechers

Outline

- How do Information-theoretic and CompMech-theoretic quantities generalize in spatially extended systems?
 - An overview of the strictly-spatial case
- How spatial Information evolves in time
 - Towards a Computational Mechanics for information processing
- The big picture: How space and time are intricately connected, spacetime probability densities, and playing outside of your light-cone.

$$H(L) = H(P(s^{L})) = -\sum_{s^{L} \in A} P(s^{L}) \log_{2} P(s^{L})$$

•What does this mean?

•What is the appropriate L for a system?

•What is s^L

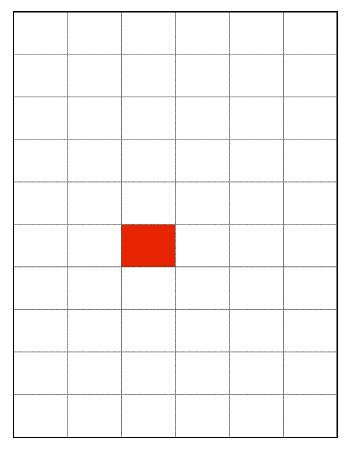
•What is the alphabet?

•How are the probabilities calculated?

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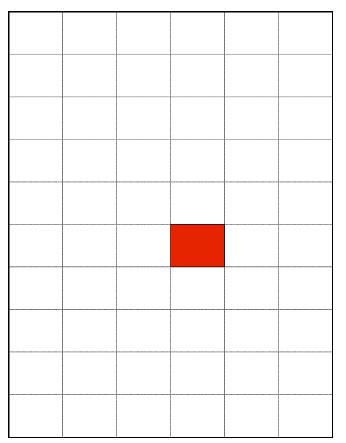
- In 1D... options
 - Consider the case of general dimensions (or at least 2D...)
 - What we should do for the 1D case will follow out as a special case

• L = 1 considers the value lone cells



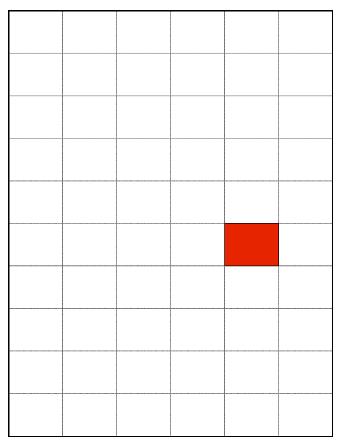
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- •But P(s) is taken from the statistics of the whole array
 - So, scan this template across the whole array to find P(s)!



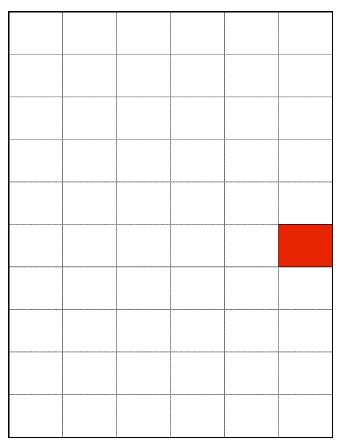
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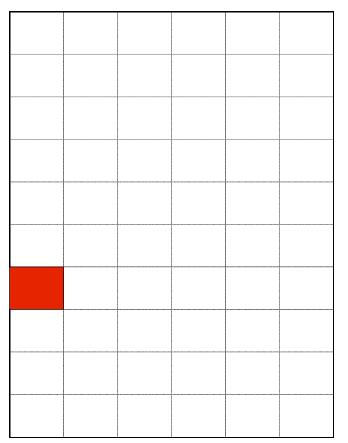
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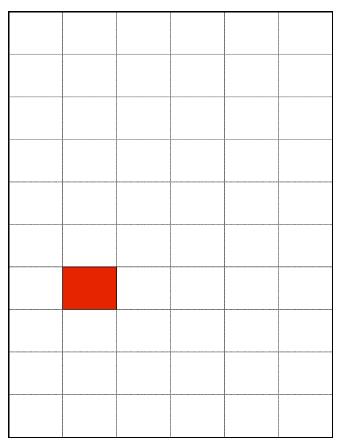
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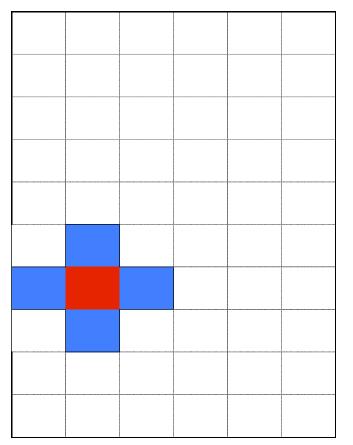
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$$H(L) = H(P(s^{L})) = -\sum_{s^{L} \in A} P(s^{L}) \log_{2} P(s^{L})$$

- L = 2 template
 depends on geometry
 of coupling (assumed
 or explicit)
- The case shown is for four nearest neighbors



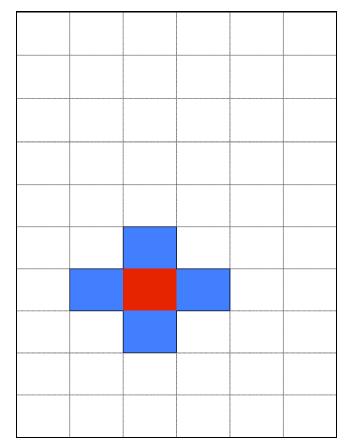
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• L = 2 template depends on geometry of coupling (assumed or explicit)

•But again, $P(s_1s_2)$ is taken from the statistics of the whole array, so we scan this template across the whole array to find $P(s_1,s_2)$.

 Notice that s_i is no longer binary!

 $s_1 \in \{0,1\}$ but $s_2 \in \{0,1,2,\dots,2^4 - 1\}$

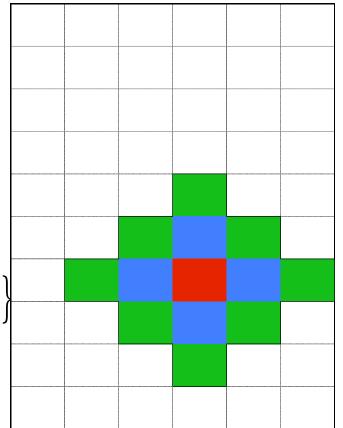


$$H(L) = H(P(s^{L})) = -\sum_{s^{L} \in A} P(s^{L}) \log_{2} P(s^{L})$$

- L = 3 template expands naturally from previous geometry
- We scan this template across the whole array to find $P(s_1, s_2, s_3)$.
- We have a different alphabet for different L. In general:

$$length\{A(L_{L^{1}-coupling})\} = \prod_{k=0}^{L-1} \left\{ n^{(\delta_{k,0}+2^{D}k^{(D-1)})} \right\}$$

for number of intrinsic states per cell, n, and coupling number (number of nearest neighbors), $\rm N_{\rm C.}$

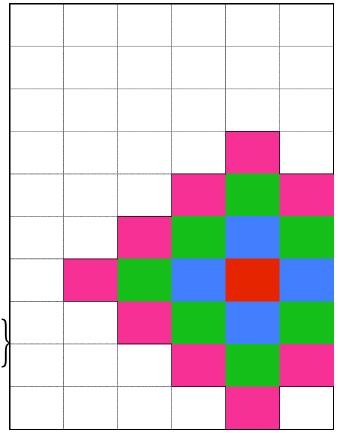


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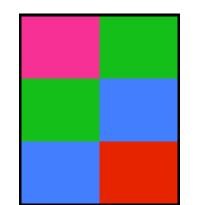
• We scan every L-template across the whole array to find $P(s^L)$.

•Special considerations for boundary conditions?

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 - H(L) saturates at L = n+m-1
 for n × m array

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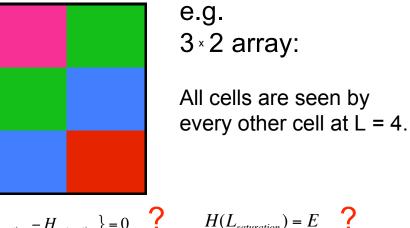
All cells are seen by every other cell at L = 4.

$$h_{\mu} = \lim_{L \to \infty} \left\{ h_{\mu}(L) \right\} = \lim_{L \to \infty} \left\{ H(L) - H(L-1) \right\} = \lim_{L \to \infty} \left\{ H_{saturation} - H_{saturation} \right\} = 0$$
?

 $H(L_{saturation}) = E \quad ?$

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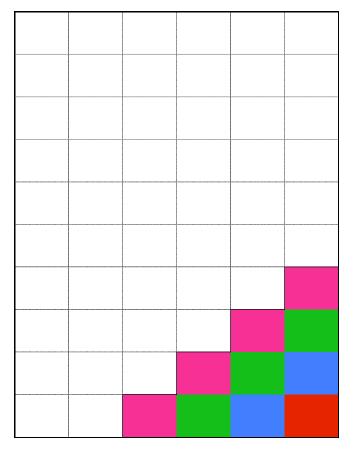
Instead, define new quantities:

$$h_{\mu}^{sat} \equiv h_{\mu}(L_{sat}) = H(L_{sat}) - H(L_{sat} - 1)$$
 ? $E^{sat} \equiv ?$

• We scan every L-template across the whole array to find $P(s^{L})$.

•Special considerations for boundary conditions?

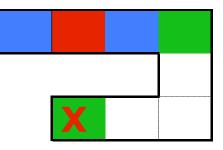
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 for n × m array
- •Just average over possibilities?



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•Special considerations for boundary conditions?

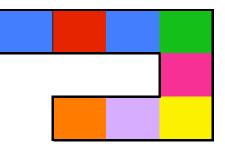
Just average over possibilities? Maybe.
Although care must be given to pathological topologies...



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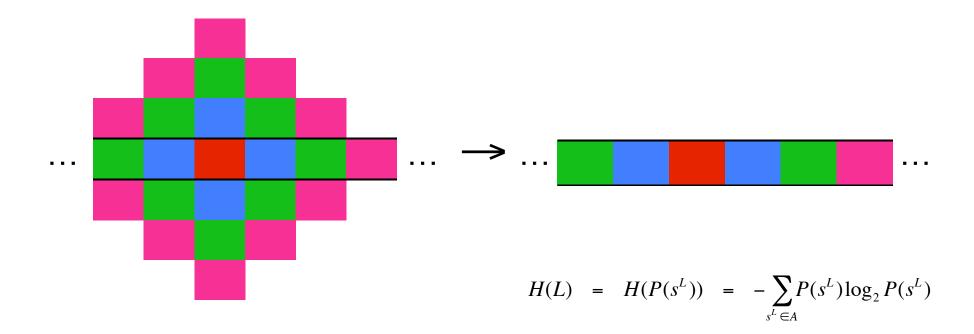
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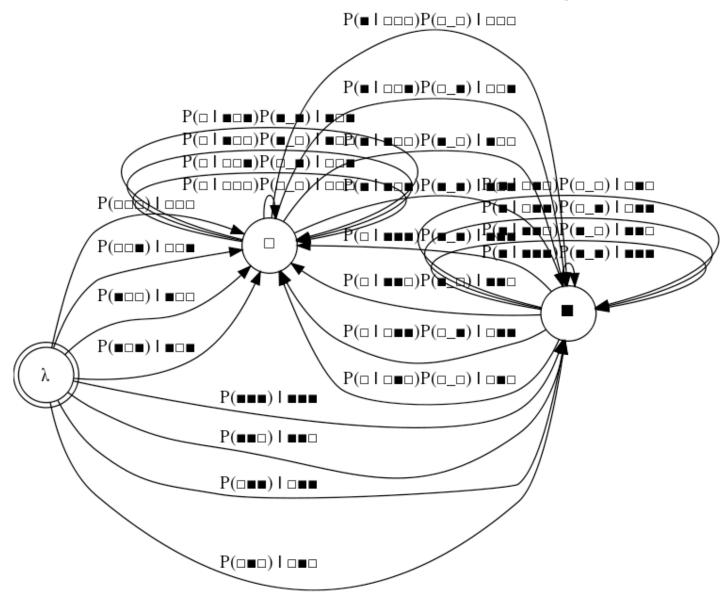
1-Dimension as boundary condition

• We scan the natural L-templates across the whole linear array to find $P(s^{L})$.

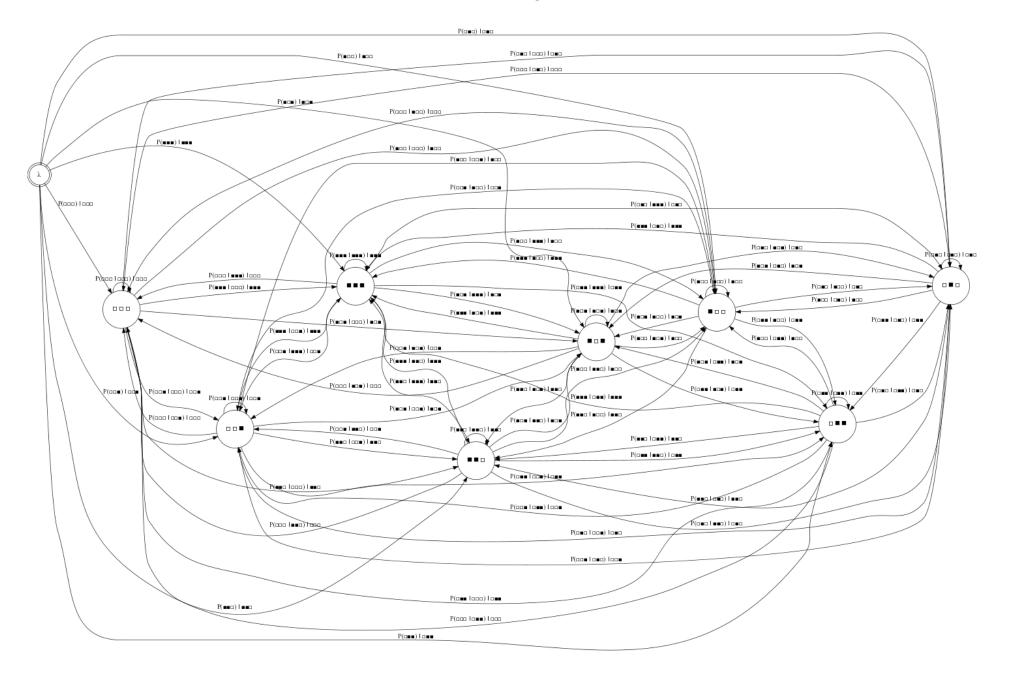


Rethinking Temporal Measurements for Spatially

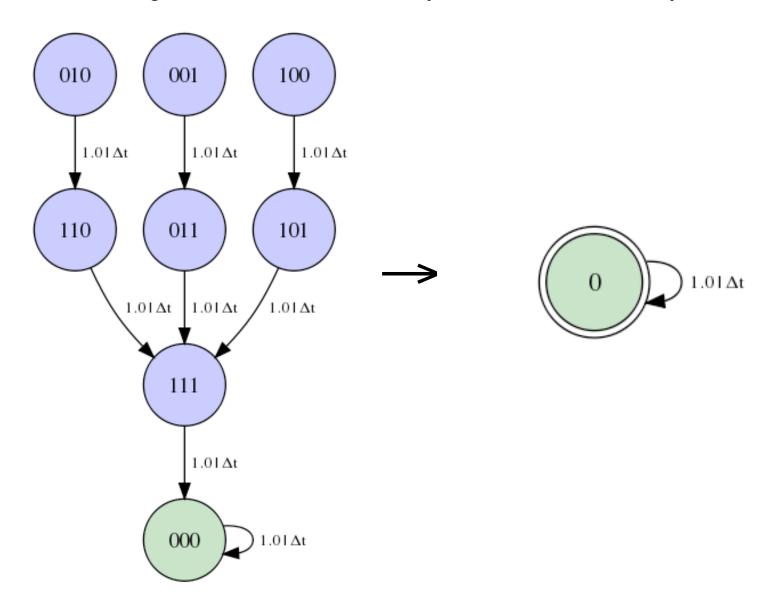
Extended States: 1-D nearest neighbor CA



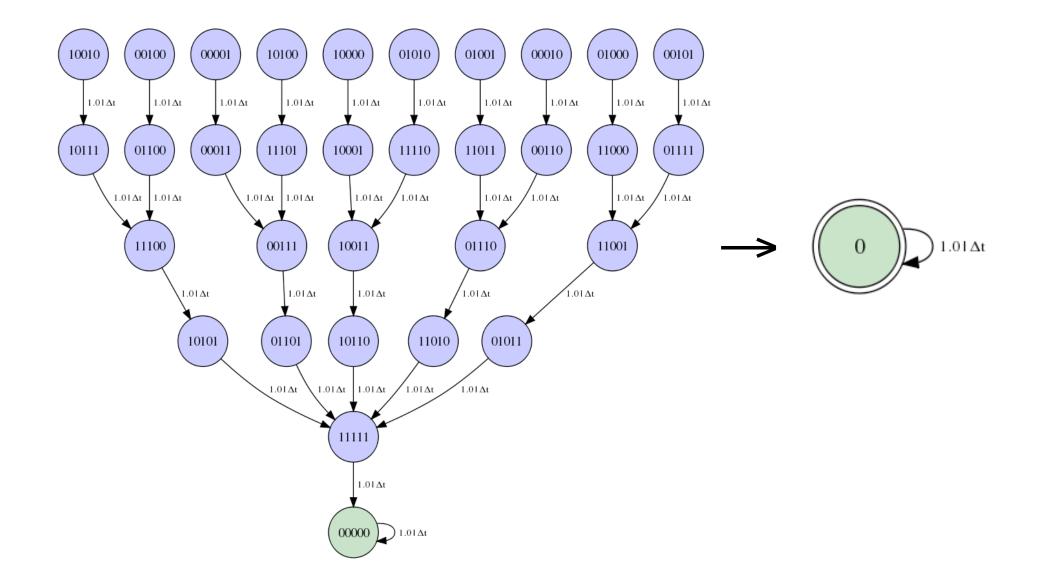
Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA



Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA Starting Small: Rule 110 with only three cells in the array



Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA Starting Small: Rule 110 with only five cells in the array



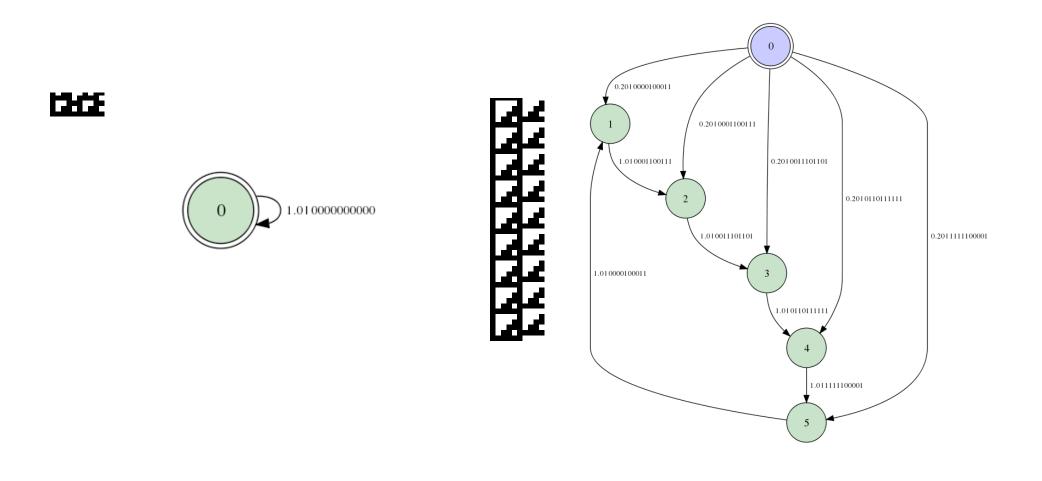
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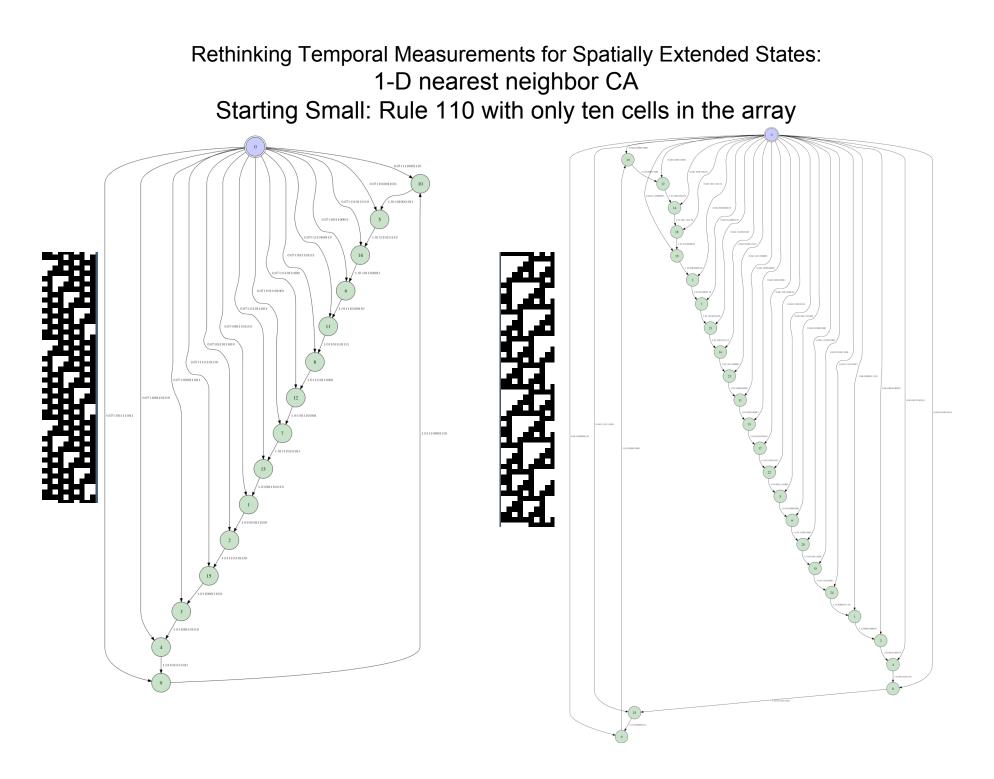


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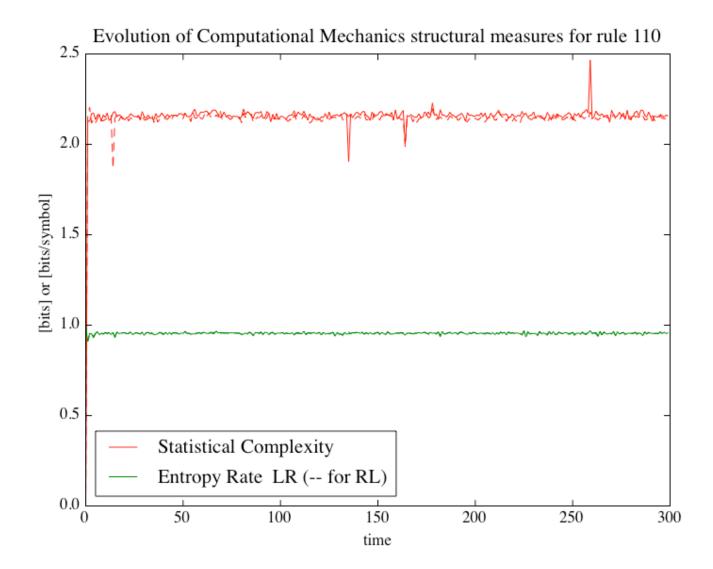
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Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA Looking Ahead: Rule 110 Impulse excitation for 3000 sites: CompMech quantities from a timeseries

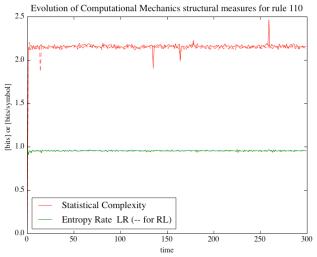


Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA Looking Ahead: Rule 110 Impulse excitation for 3000 sites: CompMech quantities from a timeseries

• The current methods/ algorithms are limited... working on that

•There are many new questions to address

• Space and time probabilities are intricately connected



Rethinking Temporal Measurements for Spatially Extended States: 1-D nearest neighbor CA

- Some inspired new measures:
 - Effective coupling length
 Helps to digitize rules of dynamic systems
 - Exploring chiral rules

