# Bidirectional Computational Mechanics II

Reading for this lecture: CMR articles TBA PRATISP IACP IACPLCOCS

# Bidirectional Computational Mechanics

Bidirectionality Excess entropy from E-machine Information diagrams for processes Bidirectional machines Bidirectional complexities

Causal shielding for forward and reverse states:

$$\Pr(\overleftarrow{X}, \overrightarrow{X} | \mathcal{S}^+) = \Pr(\overleftarrow{X} | \mathcal{S}^+) \Pr(\overrightarrow{X} | \mathcal{S}^+)$$
$$\Pr(\overleftarrow{Y}, \overrightarrow{Y} | \mathcal{S}^-) = \Pr(\overleftarrow{Y} | \mathcal{S}^-) \Pr(\overrightarrow{Y} | \mathcal{S}^-)$$

Both forward and reverse states equally good at shielding ... but for potentially different, though, related processes:

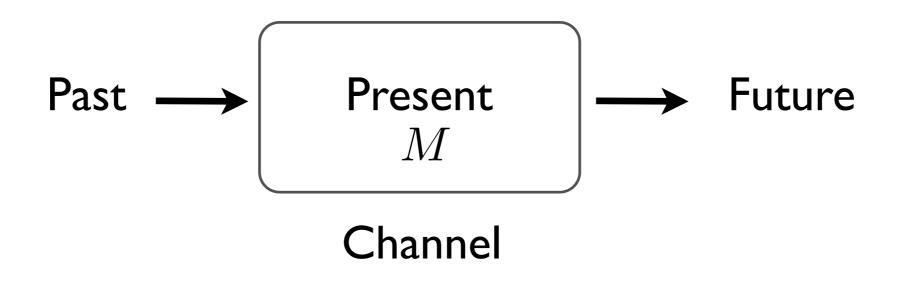
$$\overleftrightarrow{Y} = \overleftrightarrow{X}$$

## Direct relationship between forward and reverse states?

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Excess entropy from ε-machine:

Process 
$$\Pr(\overleftarrow{X}, \overrightarrow{X})$$
 is a communication channel from the past  $\overrightarrow{X}$  to the future  $\overrightarrow{X}$ :



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Excess entropy from E-machine ...

Mutual information between the past and future

$$\mathbf{E} = I[\overleftarrow{X}; \overrightarrow{X}]$$

A process's channel capacity?

More like the effective channel utilization.

## Now, how to get from given E-machine?

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Excess entropy from E-machine ...

Theorem:  $\mathbf{E} = I[\mathcal{S}^+; \mathcal{S}^-]$ 

Proof sketch: 
$$\mathbf{E} = I[\overleftarrow{X}; \overrightarrow{X}]$$
  
=  $I[\epsilon^+(\overleftarrow{X}); \epsilon^-(\overrightarrow{X})]$   
=  $I[\mathcal{S}^+; \mathcal{S}^-]$ 

 $\Rightarrow$  Need  $\Pr(\mathcal{S}^+, \mathcal{S}^-)$ 

New interpretation: Effective transmission capacity of channel between forward and reverse processes.

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Mixed state presentation of time-reversed E-machine:

$$\widetilde{M}^+ = \mathcal{T}(M^+)$$

$$M^- = \mathcal{U}(\widetilde{M}^+) \qquad \text{(minimize!)}$$

yields conditional entropy btw forward and reverse states:

 $\Pr(\mathcal{S}^+|\mathcal{S}^-)$ 

Switching maps between forward and reverse causal states:

Forward-state simplex:  $\Delta^m$   $m = |\mathcal{S}^+|$  $\Pr(\mathcal{S}_0^+ = \sigma_0, \mathcal{S}_1^+ = \sigma_1, ...) \in \Delta^m$ 

Reverse-state simplex:  $\Delta^n \quad n = |\mathcal{S}^-|$ 

$$\Pr(\mathcal{S}_0^- = \sigma_0, \mathcal{S}_1^- = \sigma_1, \ldots) \in \Delta^n$$

Forward map: 
$$f: \Delta^n \to \Delta^m$$
  
 $f(\sigma^-) = \Pr(\mathcal{S}^+ | \sigma^-)$ 

**Reverse map:**  $r: \Delta^m \to \Delta^n$ 

$$r(\sigma^+) = \Pr(\mathcal{S}^- | \sigma^+)$$

Switching maps ...

Uses:

Calculate with conditional & joint state dependencies:

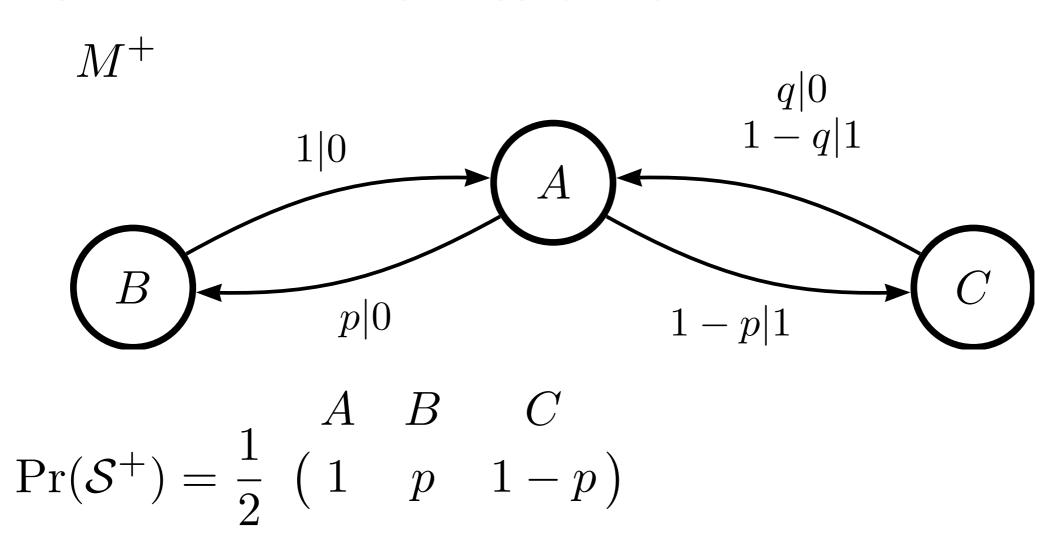
 $\Pr(\mathcal{S}^+|\mathcal{S}^-)$  $\Pr(\mathcal{S}^+,\mathcal{S}^-)$ 

Recast quantifiers purely in terms of forward & reverse states.

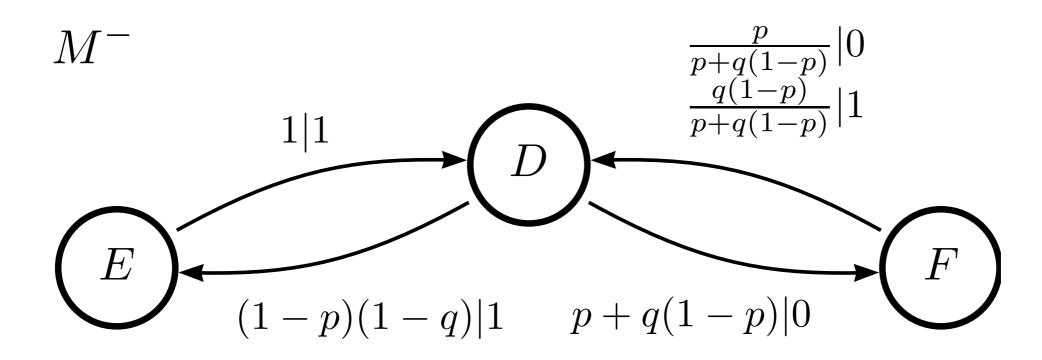
New quantifiers.

A new representation ...

## Example: Random Noisy Copy (RnC)



Example: Random Noisy Copy (RnC)



$$\Pr(\mathcal{S}^{-}) = \frac{1}{2} \begin{pmatrix} D & E & F \\ 1 & (1-p)(1-q) & p+q(1-p) \end{pmatrix}$$

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Example: Random Noisy Copy (RnC)

Forward switching map: Conditional distribution, from previous MSP calculation

$$\Pr(\mathcal{S}^{+}|\mathcal{S}^{-}) = \begin{bmatrix} A & B & C \\ D & 0 & 1 \\ E & 1 & 0 & 0 \\ F & 0 & \frac{p}{p+q(1-p)} & \frac{q(1-p)}{p+q(1-p)} \end{bmatrix}$$

## Example: Random Noisy Copy (RnC)

Joint distribution:

$$\Pr(\mathcal{S}^+, \mathcal{S}^-) = \frac{1}{2} \times \begin{bmatrix} D \\ E \\ F \end{bmatrix} \begin{pmatrix} 0 & 0 & 1 \\ (1-p)(1-q) & 0 & 0 \\ 0 & p & q(1-p) \end{pmatrix}$$

Componentwise calculation:

$$\Pr(\mathcal{S}^+ = i, \mathcal{S}^- = j) = \Pr(\mathcal{S}^+ = i|\mathcal{S}^- = j)\Pr(\mathcal{S}^- = j)$$

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Example: Random Noisy Copy (RnC)

Reverse switching map:

"Transpose" of joint distribution

$$\Pr(\mathcal{S}^{-}, \mathcal{S}^{+}) = \frac{1}{2} \times \begin{pmatrix} A \\ B \\ C \end{pmatrix} \begin{pmatrix} 0 & (1-p)(1-q) & 0 \\ 0 & 0 & p \\ 1 & 0 & q(1-p) \end{pmatrix}$$

/ )

F

F

Componentwise:

$$\Pr(\mathcal{S}^{-} = i | \mathcal{S}^{+} = j) = \frac{\Pr(\mathcal{S}^{-} = i, \mathcal{S}^{+} = j)}{\Pr(\mathcal{S}^{+} = j)}$$
$$D \qquad E \qquad F$$
$$\Pr(\mathcal{S}^{-} | \mathcal{S}^{+}) = \frac{A}{B} \begin{pmatrix} 0 & (1-p)(1-q) & 0\\ 0 & 0 & 1\\ \frac{1}{1-p} & 0 & q \end{pmatrix}$$

Example: Random Noisy Copy (RnC)

Reverse switching map ...

Normalize rows to get actual transition probabilities:

$$\Pr(\mathcal{S}^{-}|\mathcal{S}^{+}) = \begin{array}{ccc} A \\ B \\ C \end{array} \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ \frac{1}{1+q(1-p)} & 0 & \frac{q(1-p)}{1+q(1-p)} \end{pmatrix}$$

## Example: Random Noisy Copy (RnC) ...

## Excess entropy:

$$\mathbf{E} = I[\mathcal{S}^{+}; \mathcal{S}^{-}] = H[\mathcal{S}^{+}] - H[\mathcal{S}^{+}|\mathcal{S}^{-}] = H\left[\frac{1}{2}(1 \ p \ 1-p)\right] - H[\{B, C\}|F] \Pr(F) = H\left[\frac{1}{2}(1 \ p \ 1-p)\right] - H\left(\frac{p}{p+q(1-p)}\right)(p+q(1-p))$$

$$\Pr(\mathcal{S}^+) = \frac{1}{2} \begin{pmatrix} A & B & C \\ (1 & p & 1-p \end{pmatrix} & A & B & C \\ & & & \\ & & & \\ & &$$

Crypticities, recast: Forward:

$$\chi^{+} = H[\mathcal{S}^{+}|\overrightarrow{X}]$$
$$= H[\mathcal{S}^{+}|\epsilon^{-}(\overrightarrow{X})]$$
$$= H[\mathcal{S}^{+}|\mathcal{S}^{-}]$$

#### Reverse:

$$\chi^{-} = H[\mathcal{S}^{-}|\overleftarrow{X}]$$
$$= H[\mathcal{S}^{-}|\epsilon^{+}(\overleftarrow{X})]$$
$$= H[\mathcal{S}^{-}|\mathcal{S}^{+}]$$

Example: Random Noisy Copy (RnC) ...

$$\chi^{+} = H[\mathcal{S}^{+}|\mathcal{S}^{-}]$$
$$= H[\mathcal{S}^{+} = \{B, C\}|\mathcal{S}^{-} = F]$$
$$= H\left(\frac{p}{p+q(1-p)}\right)(p+q(1-p))$$

$$\chi^{-} = H[\mathcal{S}^{-}|\mathcal{S}^{+}]$$
$$= H\left[\mathcal{S}^{-} = \{D, F\} | \mathcal{S}^{+} = C\right]$$
$$= H\left(\frac{1}{1+q(1-p)}\right)\left(\frac{1-p}{2}\right)$$

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Causal irreversibility, recast:

$$\Xi \equiv C_{\mu}^{+} - C_{\mu}^{-}$$
$$= H[\mathcal{S}^{+}|\mathcal{S}^{-}] - H[\mathcal{S}^{-}|\mathcal{S}^{+}]$$

using

$$C^+_{\mu} = \mathbf{E} + \chi^+$$
$$C^-_{\mu} = \mathbf{E} + \chi^-$$

Example: Random Noisy Copy (RnC) ...

$$\Xi = H[S^{+}|S^{-}] - H[S^{-}|S^{+}] = H\left[\frac{1}{2}(1 \ p \ 1-p)\right] - H\left[\frac{1}{2}(1 \ (1-p)(1-q) \ p+q(1-p))\right]$$

Summary:

New meaning for excess entropy:

$$\mathbf{E} = I[\mathcal{S}^+; \mathcal{S}^-]$$

 ${f E}$  can be directly (and accurately!) calculated from  $\epsilon$ -machine.

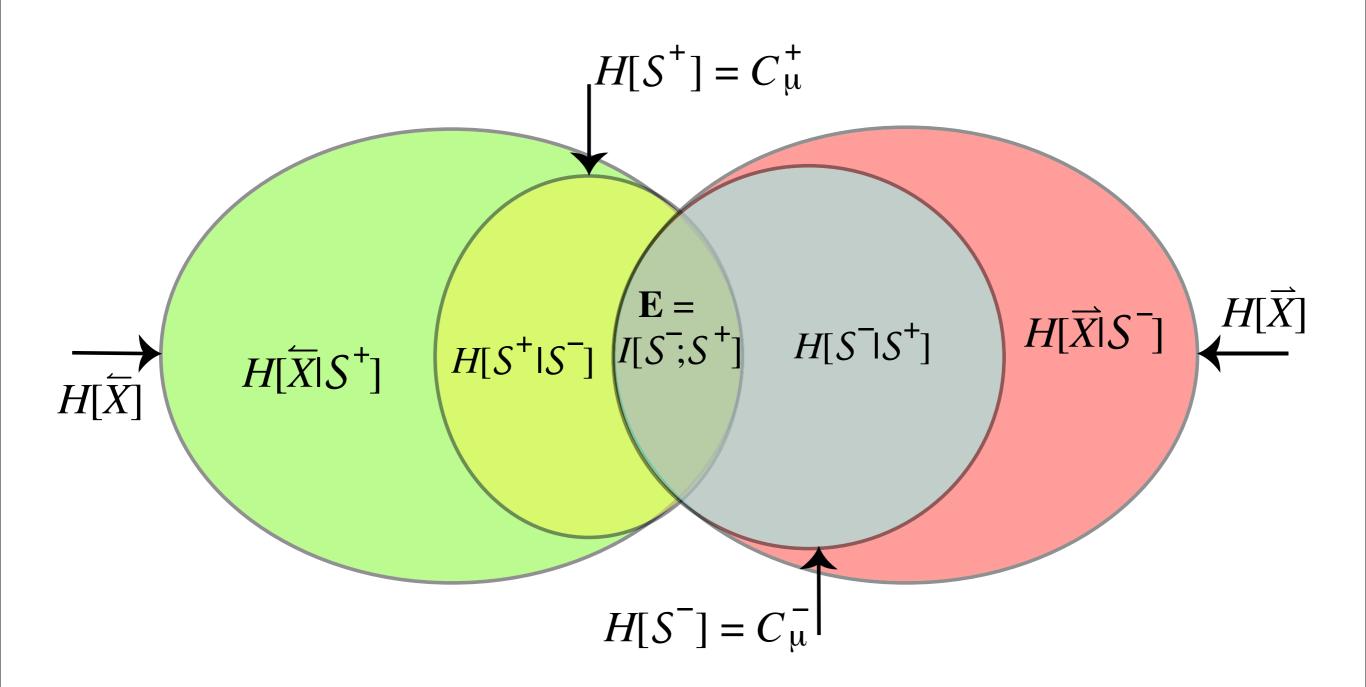
New level of analytical calculation possible.

New algorithms for measuring intrinsic computation using bidirectional representations.

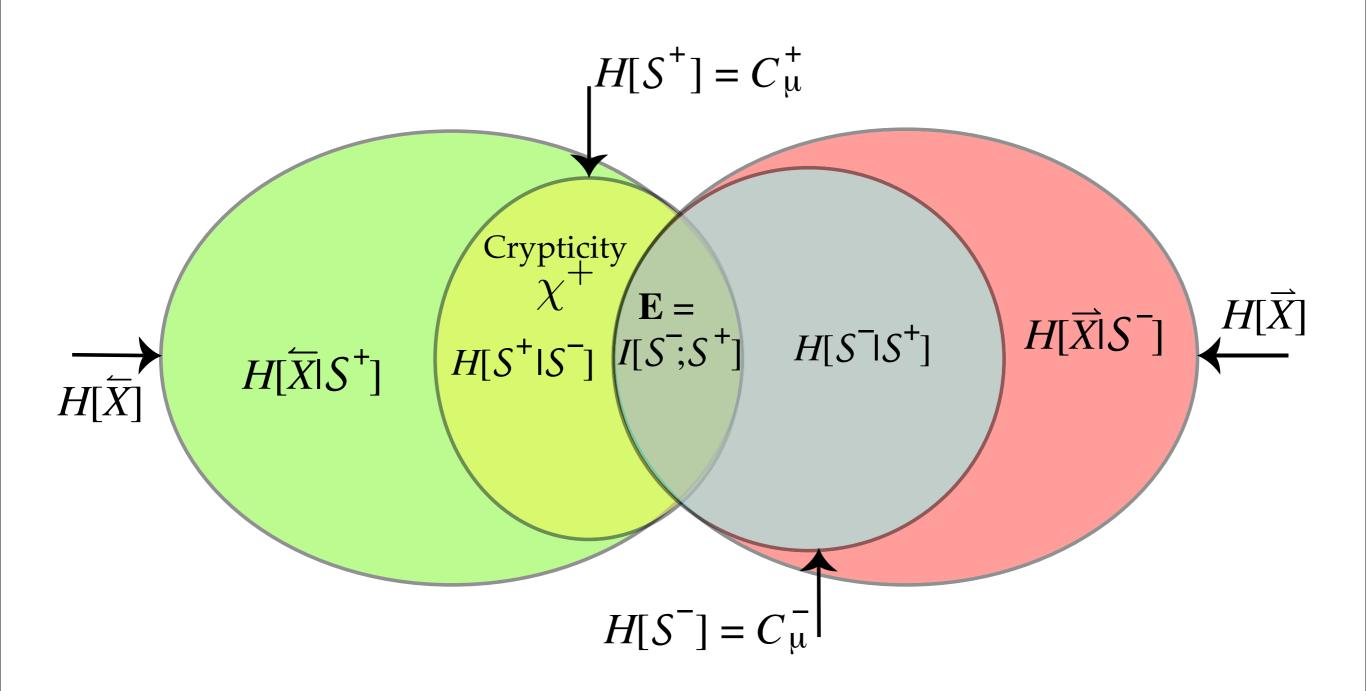
Now, ready for full information diagram of a process.

Requires both forward and reverse scanning of a process.

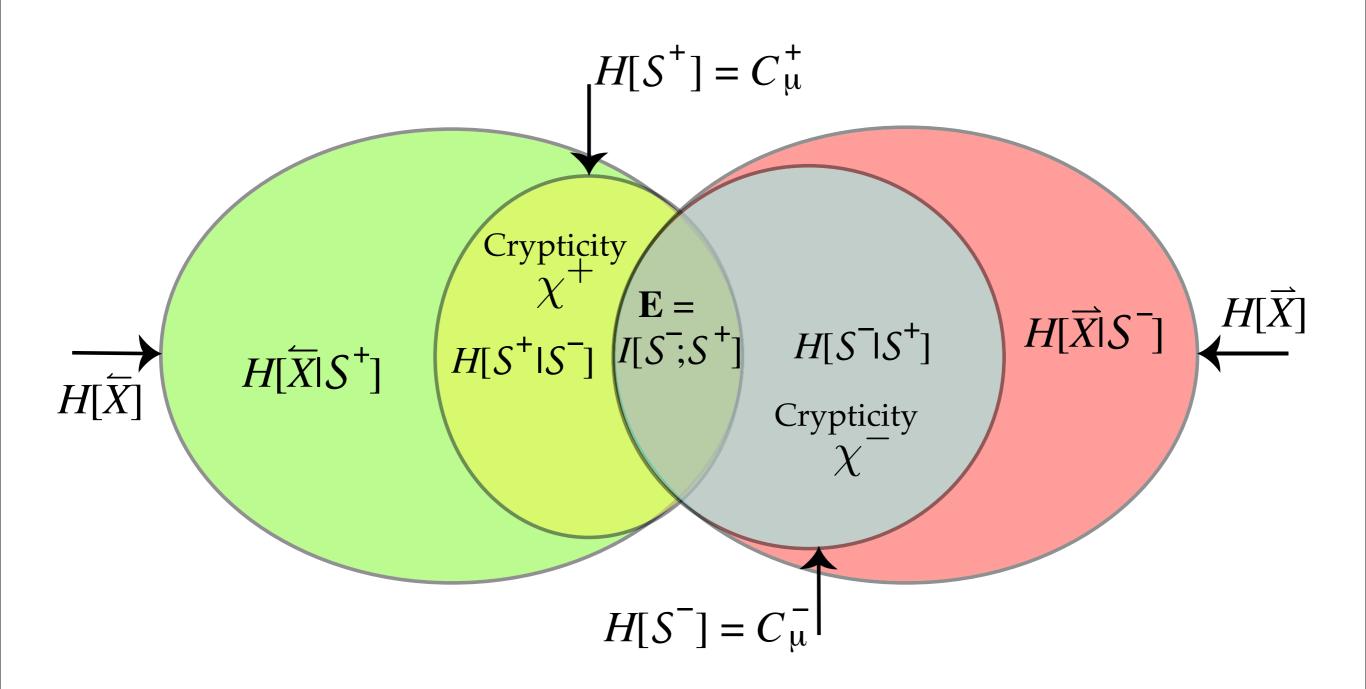
#### $\epsilon\textsc{-machine}$ information Diagram



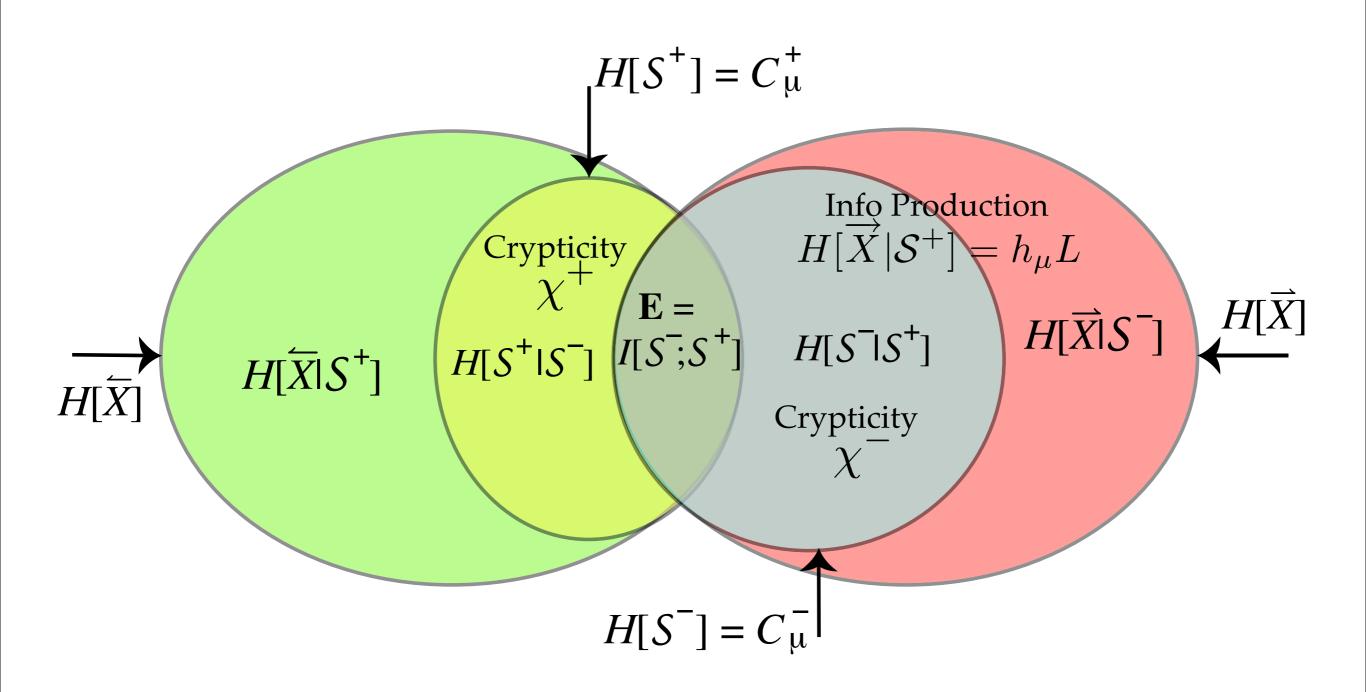
## $\epsilon\textsc{-machine}$ information Diagram



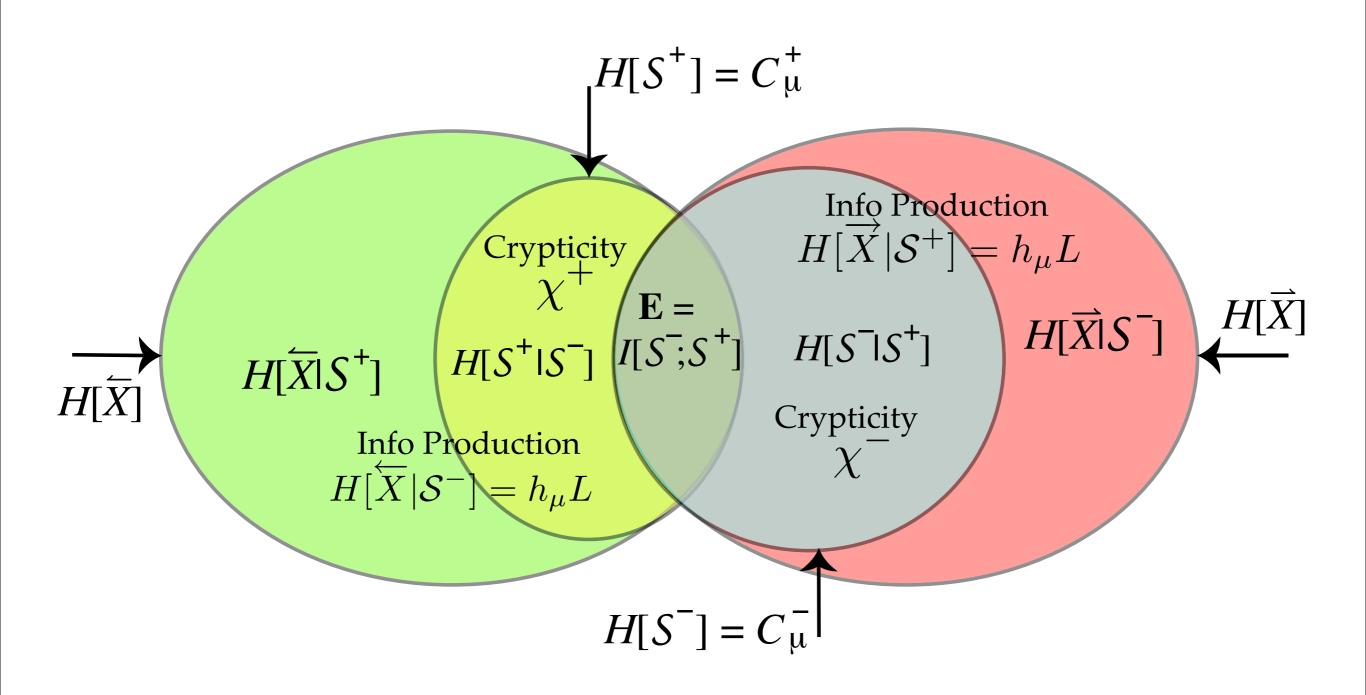
#### $\epsilon\textsc{-machine}$ information Diagram



## $\epsilon\textsc{-machine}$ information Diagram



## $\epsilon\textsc{-machine}$ information Diagram



Summary:

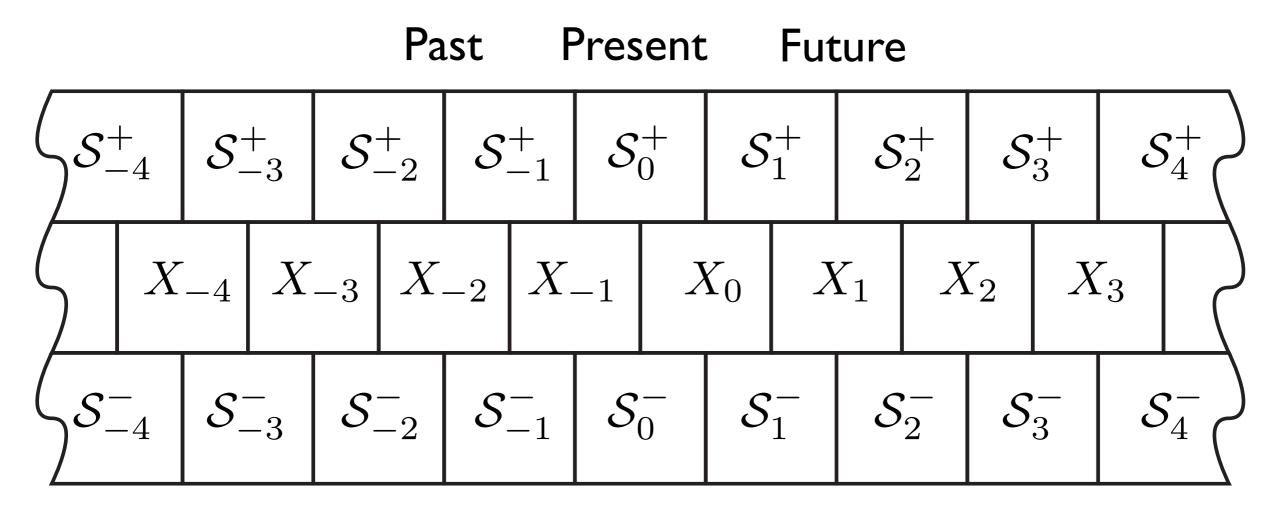
## So far, re-expressed key measures in terms of forward and reverse E-machines

$$\mathbf{E} = I[\mathcal{S}^{-}; \mathcal{S}^{+}]$$
$$\chi^{+} = H[\mathcal{S}^{+}|\mathcal{S}^{-}]$$
$$\chi^{-} = H[\mathcal{S}^{-}|\mathcal{S}^{+}]$$
$$\Xi = H[\mathcal{S}^{+}|\mathcal{S}^{-}] - H[\mathcal{S}^{-}|\mathcal{S}^{+}]$$

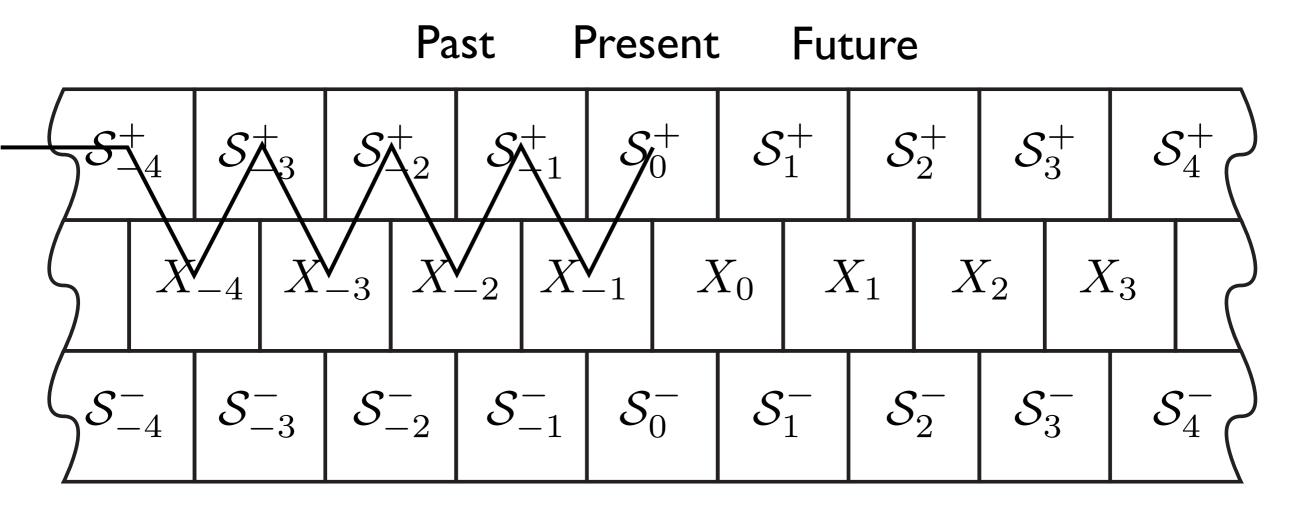
# How to more directly represent the interdependence between forward and reverse E-machines?

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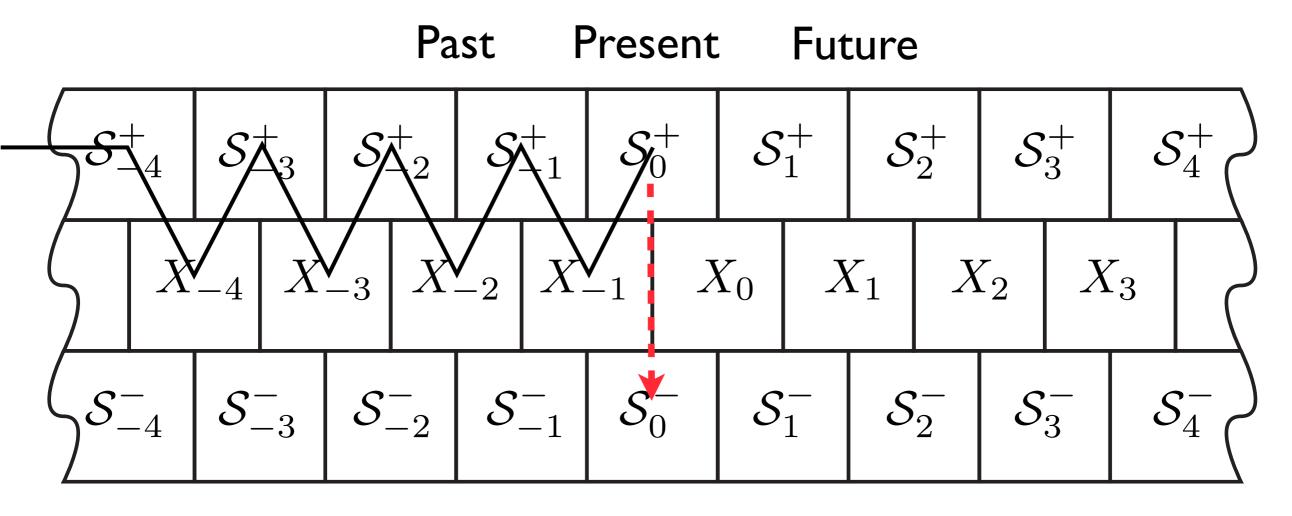
#### **Bidirectional Process Lattice:**



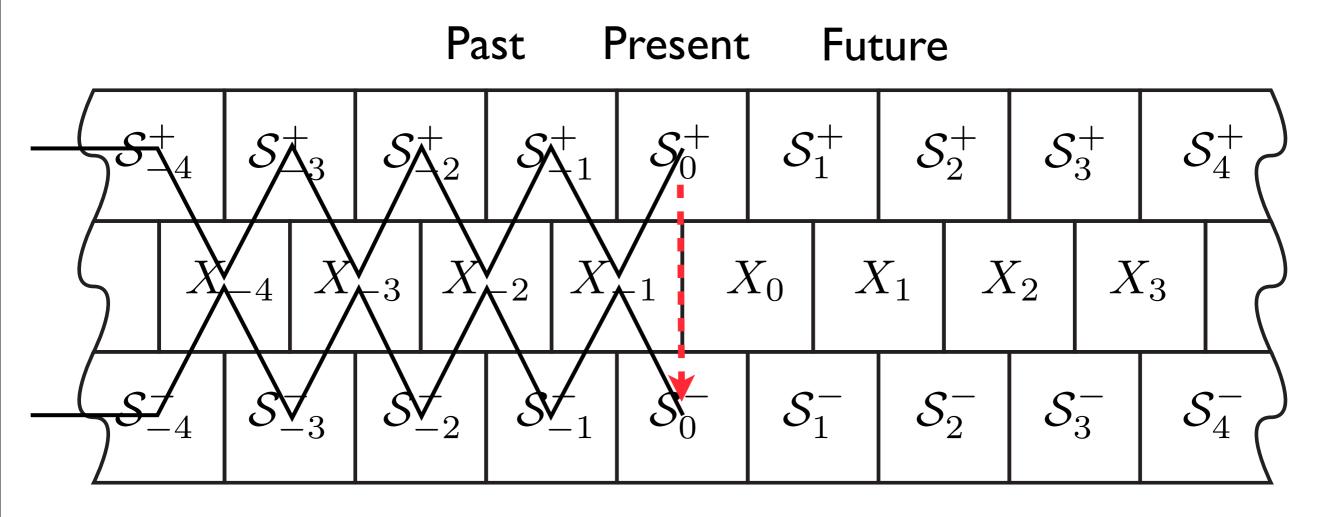
#### **Bidirectional Process Lattice:**



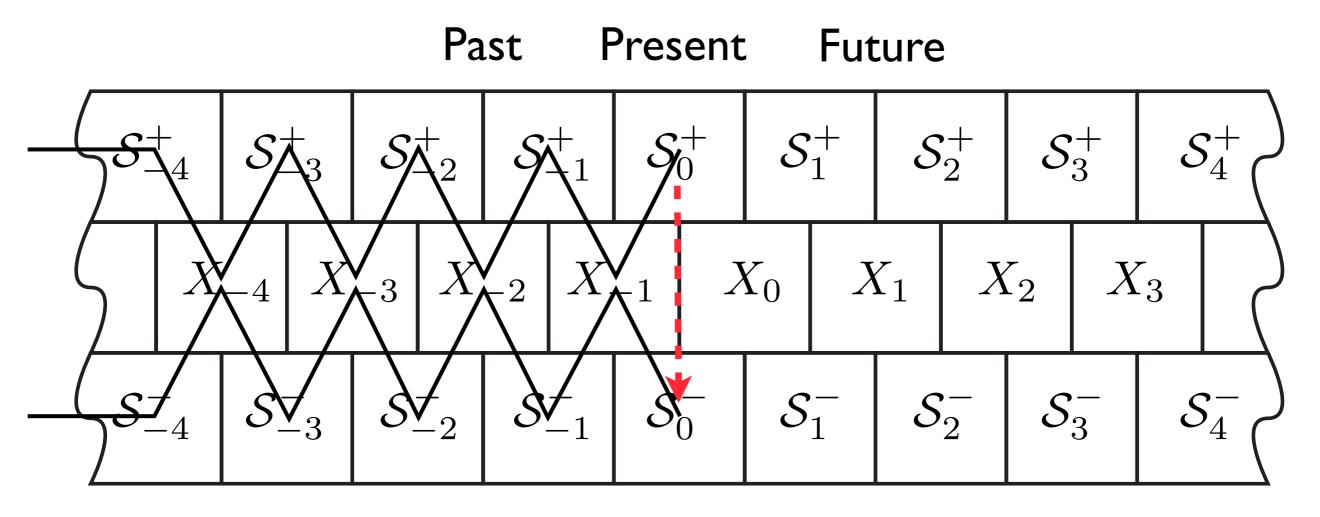
#### **Bidirectional Process Lattice:**



#### **Bidirectional Process Lattice:**



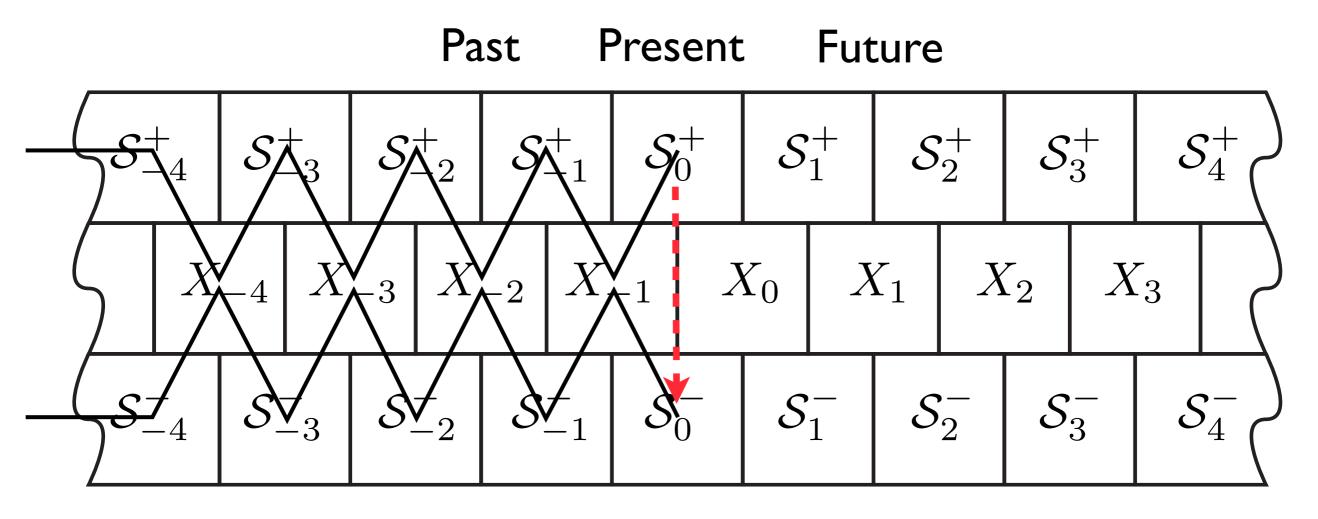
## **Bidirectional Process Lattice:**



You can choose to reverse direction. Path automaton over scan direction  $\{+, -\}$ .

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# **Bidirectional Process Lattice:**



You can choose to reverse direction. Path automaton over scan direction  $\{+, -\}$ .

What describes this?

A bidirectional machine, with forward and reverse moves. Previous forward & reverse machines are a subset of paths.

Bidirectional machine:  $M^{\pm}$ 

Equivalence relation:  $\sim^{\pm}$ 

$$\epsilon^{\pm}(\overleftarrow{x}) = \left\{ \overleftarrow{x}' = \overleftarrow{x}' \overrightarrow{x}' : \overleftarrow{x}' \in \epsilon^{+}(\overleftarrow{x}) \text{ and } \overrightarrow{x}' \in \epsilon^{-}(\overrightarrow{x}) \right\}$$

**Bidirectional states:** 

$$\begin{aligned} \boldsymbol{\mathcal{S}}^{\pm} &= \operatorname{Pr}(\overleftarrow{X}, \overrightarrow{X}) / \sim^{\pm} & \text{A partition of } \overleftarrow{X} \\ &\subseteq \boldsymbol{\mathcal{S}}^{+} \times \boldsymbol{\mathcal{S}}^{-} \end{aligned}$$

**Bidirectional Machine:** 

$$M^{\pm} = \{ \boldsymbol{\mathcal{S}}^{\pm}; \mathcal{T}^{(x)}, x \in \mathcal{A} \}$$

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The bidirectional causal state the process is in at time t is

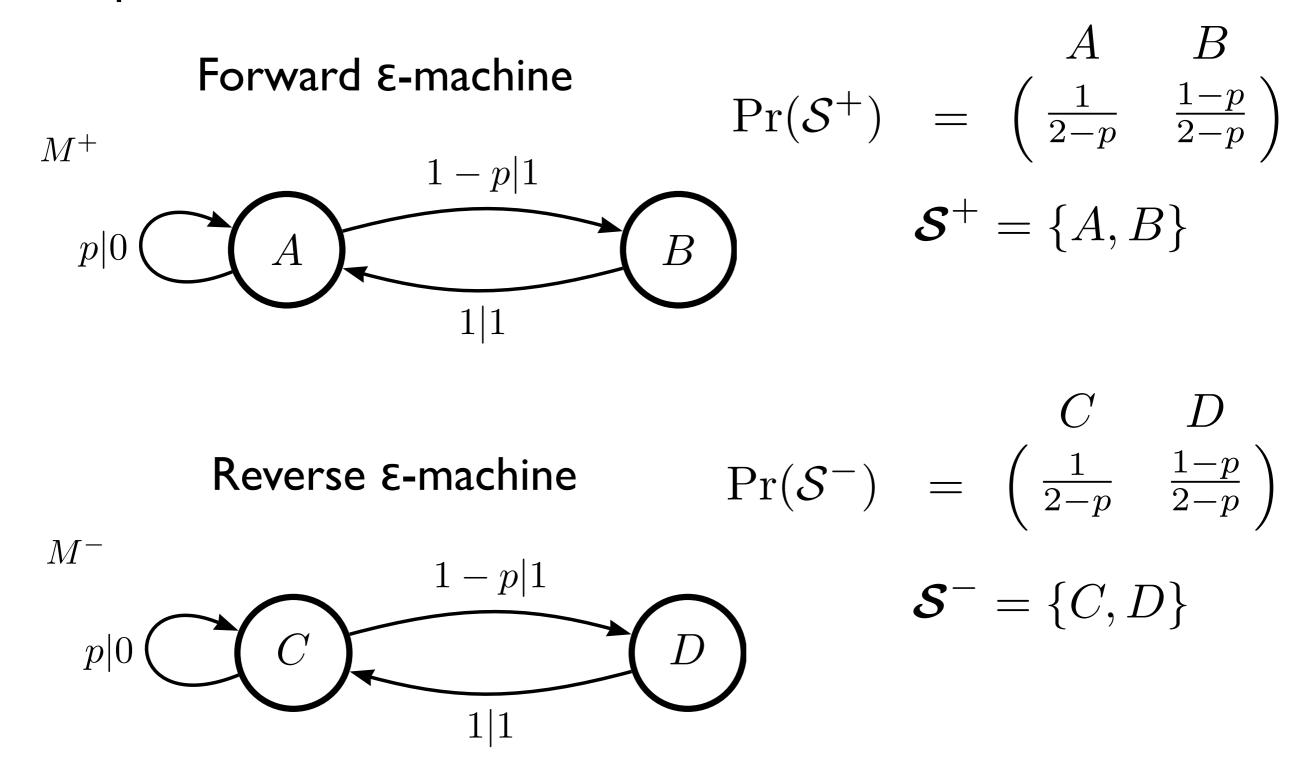
$$S_{t}^{\pm} = (\epsilon^{+}(\overleftarrow{x}_{t}), \epsilon^{-}(\overrightarrow{x}_{t})) \qquad \quad \forall \vec{x} = \overleftarrow{x}_{t} \vec{x}_{t}$$
Past Present Future
$$S_{-4}^{+} S_{-3}^{+} S_{-2}^{+} S_{-1}^{+} S_{0}^{+} S_{1}^{+} S_{2}^{+} S_{3}^{+} S_{4}^{+}$$

$$X_{-4} X_{-3} X_{-2} X_{-1} X_{0} X_{1} X_{2} X_{3}$$

$$S_{-4}^{-} S_{-3}^{-} S_{-2}^{-} S_{-1}^{-} S_{0}^{-} S_{1}^{-} S_{2}^{-} S_{3}^{-} S_{4}^{-}$$

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# **Example: Even Process**



Example: Even Process ...

$$h_{\mu} = H(p)/(2-p)$$
  
 $C_{\mu}^{+} = H(1/(2-p))$   
 $C_{\mu}^{-} = H(1/(2-p))$   
 $\Xi = 0$ 

# Example: Even Process ...

Forward switching map:

$$\Pr(\mathcal{S}^+|\mathcal{S}^-) = \frac{C}{D} \begin{pmatrix} 1 & 0\\ 0 & 1 \end{pmatrix}$$

Reverse switching map:

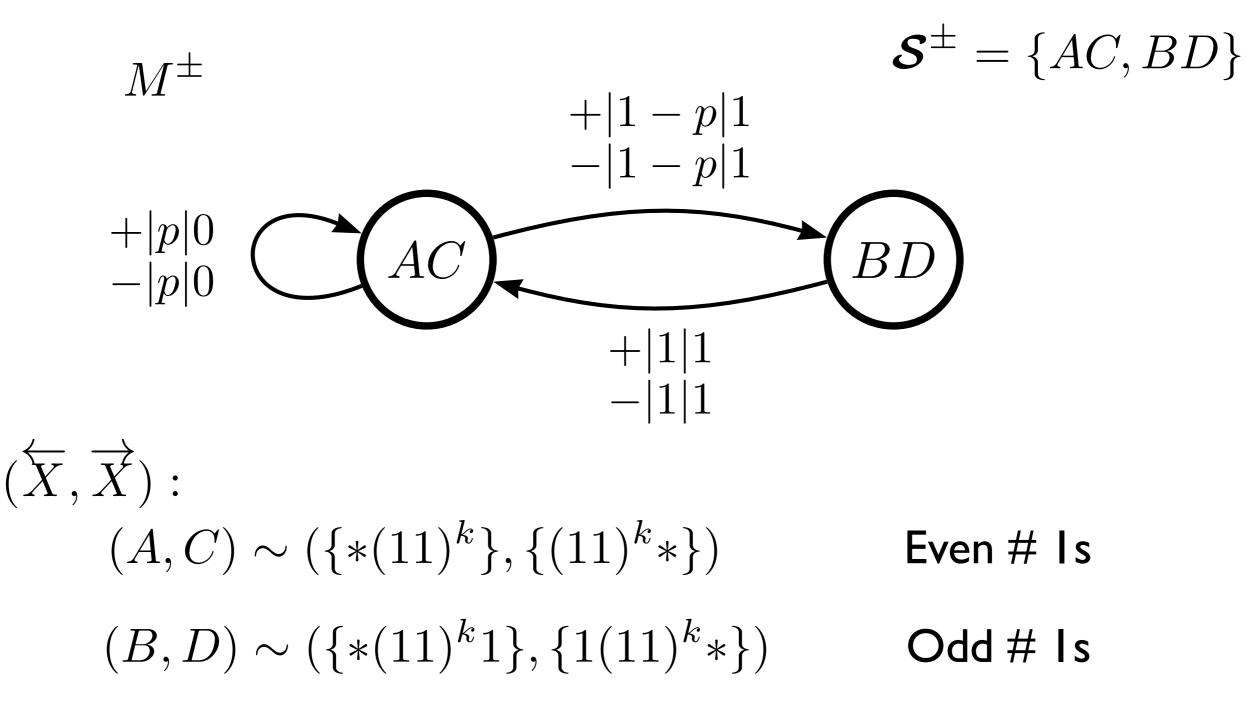
$$\Pr(\mathcal{S}^{-}|\mathcal{S}^{+}) = \frac{A}{B} \begin{pmatrix} 1 & 0\\ 0 & 1 \end{pmatrix}$$

# Identities!

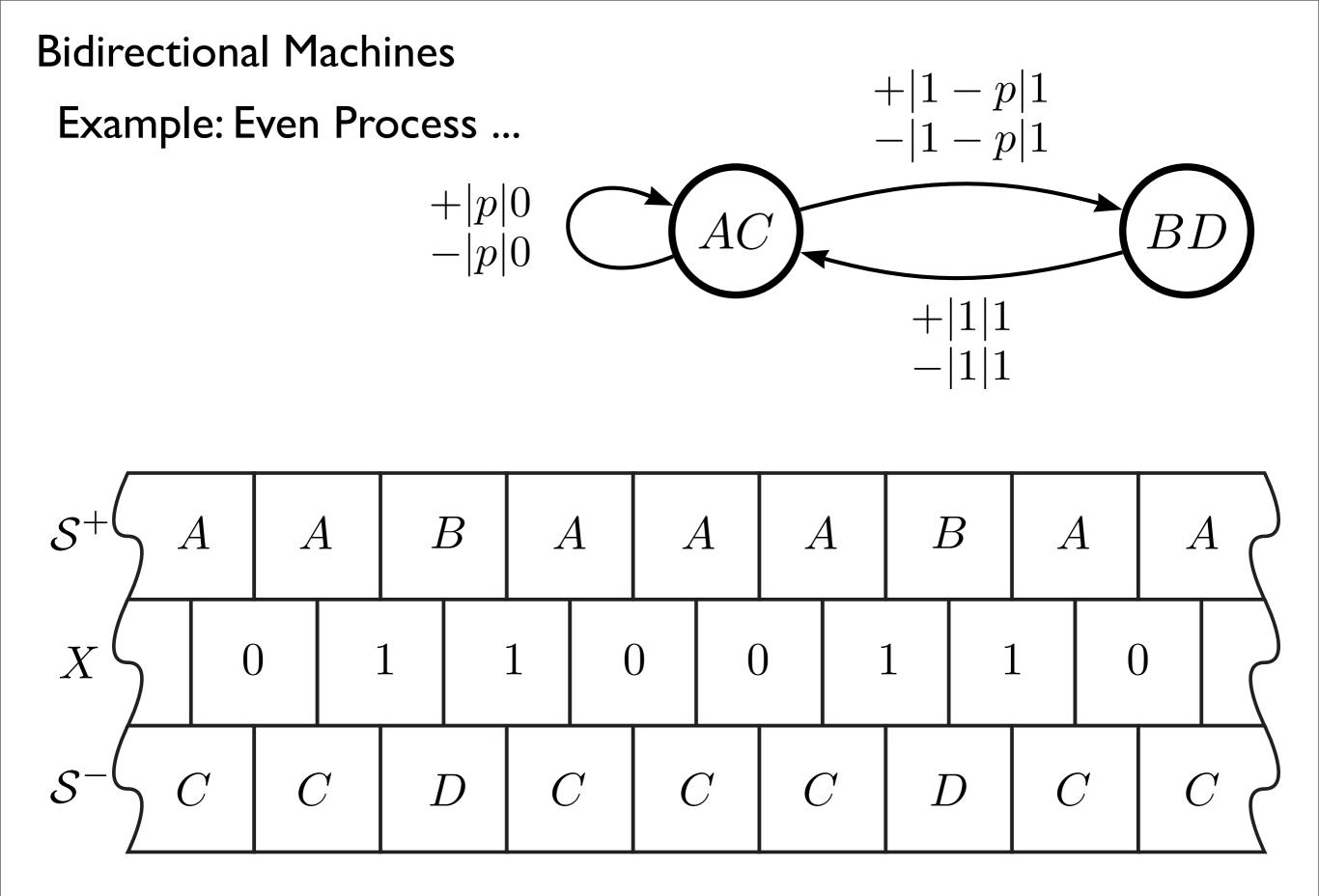
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# Example: Even Process ...

**Bidirectional machine:** 



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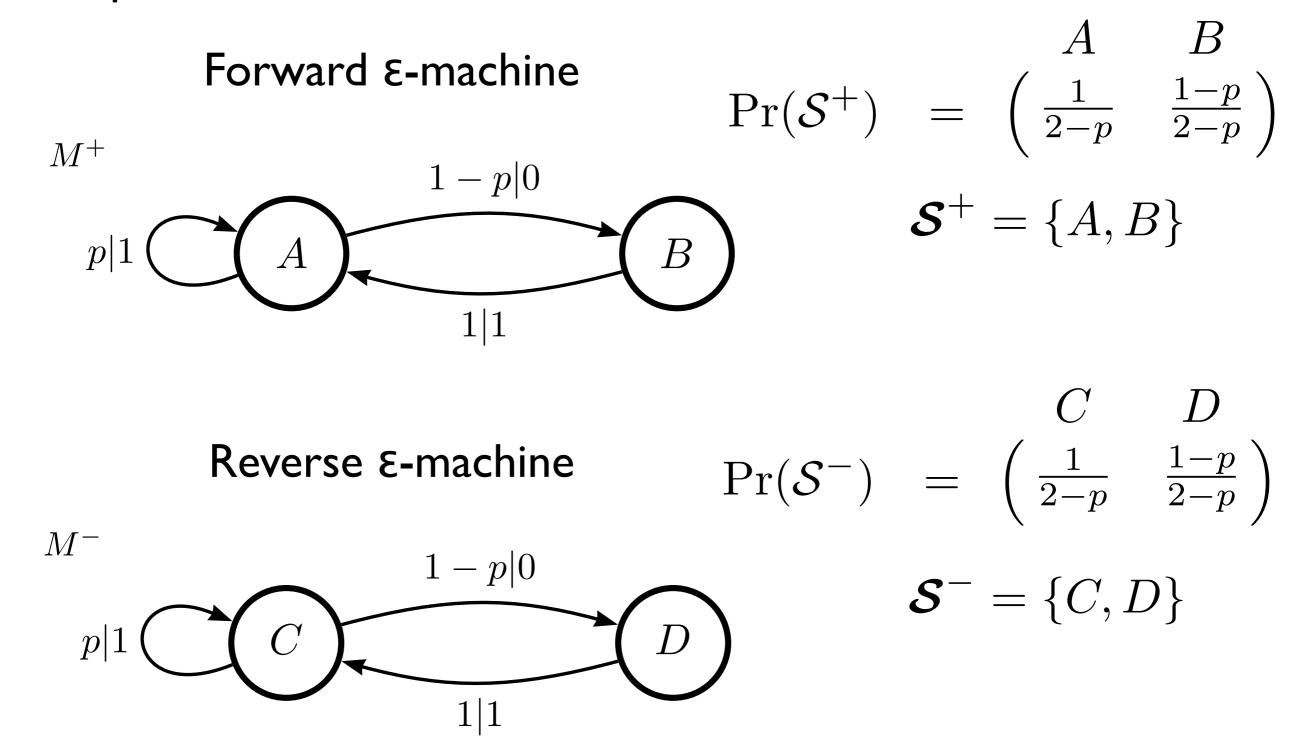
Example: Even Process ...

$$\chi^+(p) = \chi^-(p) = 0$$

$$\mathbf{E} = C_{\mu}^{+} = C_{\mu}^{-}$$

Sofic, non-Markov But simple! Microscopically reversible Causally reversible Explicit process: non-cryptic

#### **Example: Golden Mean Process**



Example: Golden Mean Process ...

$$h_{\mu} = H(p)/(2-p)$$
  
 $C_{\mu}^{+} = H(1/(2-p))$   
 $C_{\mu}^{-} = H(1/(2-p))$   
 $\Xi = 0$ 

# Same as Even Process!

Example: Golden Mean Process ...

Forward switching map:

$$\Pr(\mathcal{S}^+|\mathcal{S}^-) = \frac{C}{D} \begin{pmatrix} p & 1-p\\ 1 & 0 \end{pmatrix}$$

Reverse switching map: 
$$C$$
  $D$   
 $\Pr(\mathcal{S}^{-}|\mathcal{S}^{+}) = \frac{A}{B} \begin{pmatrix} p & 1-p \\ 1 & 0 \end{pmatrix}$ 

# Not identities!

Ambiguity and loss of information on switching.

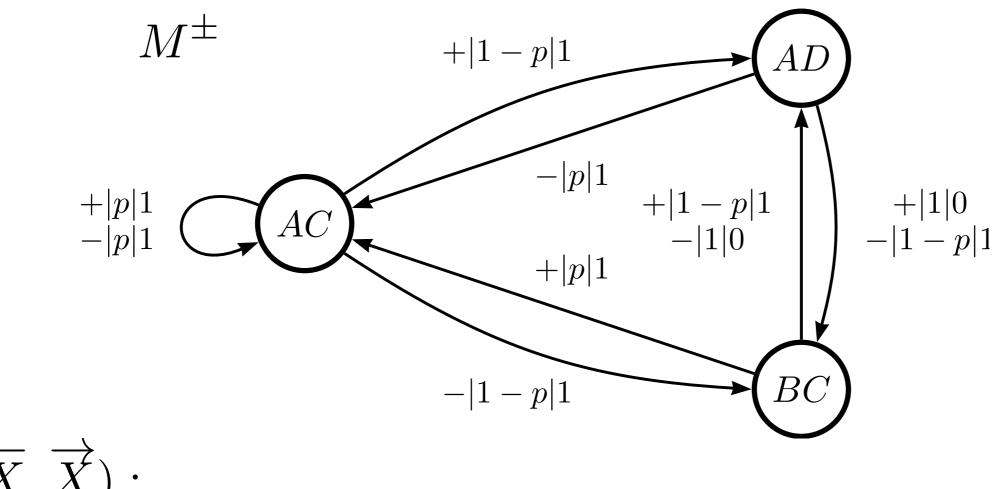
A

B

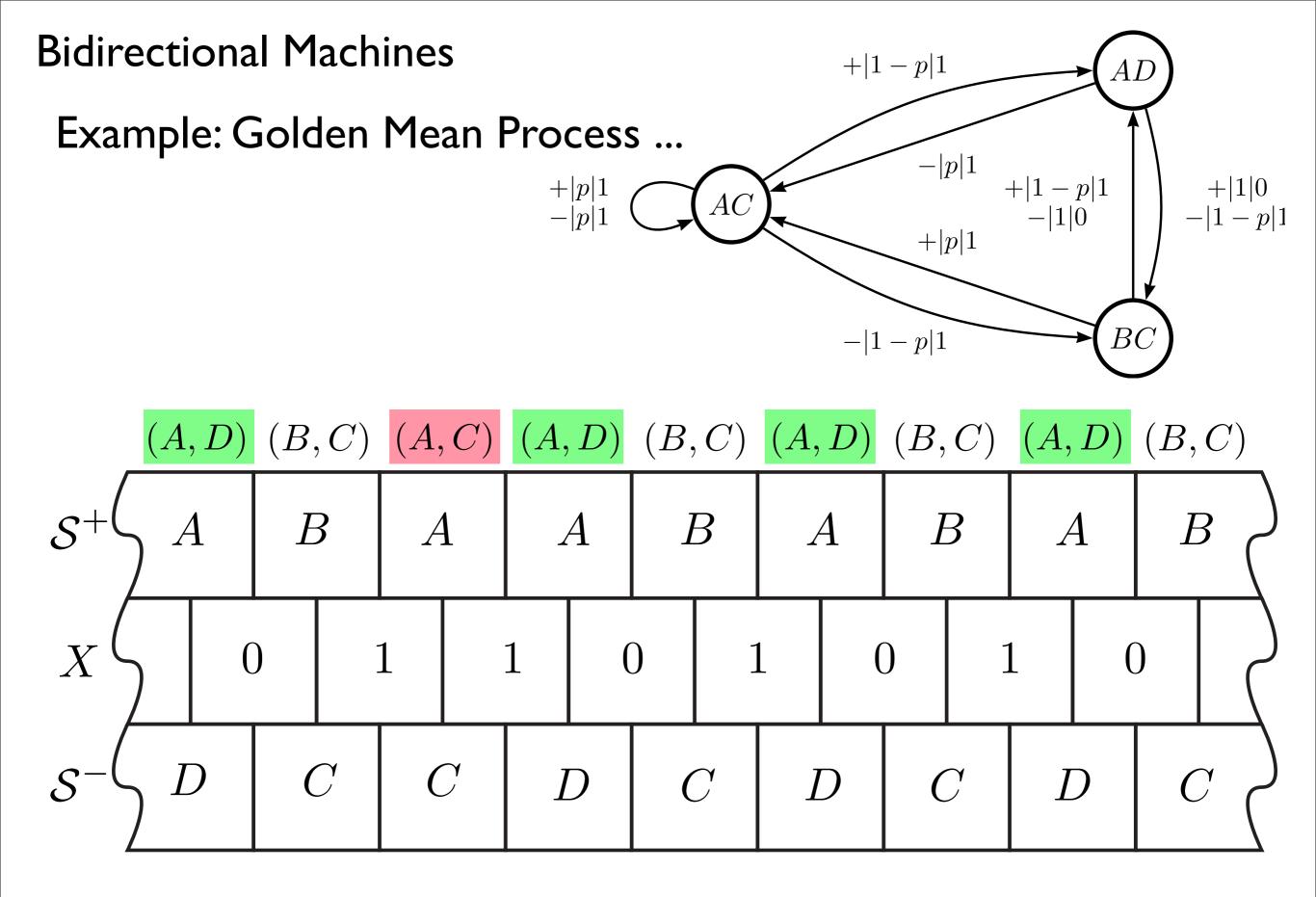
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# Example: Golden Mean Process ... Bidirectional machine:





$$(X, X): (A, C) \sim (\{*1\}, \{1*\}) (A, D) \sim (\{*1\}, \{0*\}) (B, C) \sim (\{*0\}, \{1*\})$$



Example: Golden Mean Process ...

Order-I Markov:

$$\mathbf{E} = C_{\mu}^{+} - h_{\mu}$$
$$= H\left(\frac{1}{2-p}\right) - \frac{H(p)}{2-p}$$

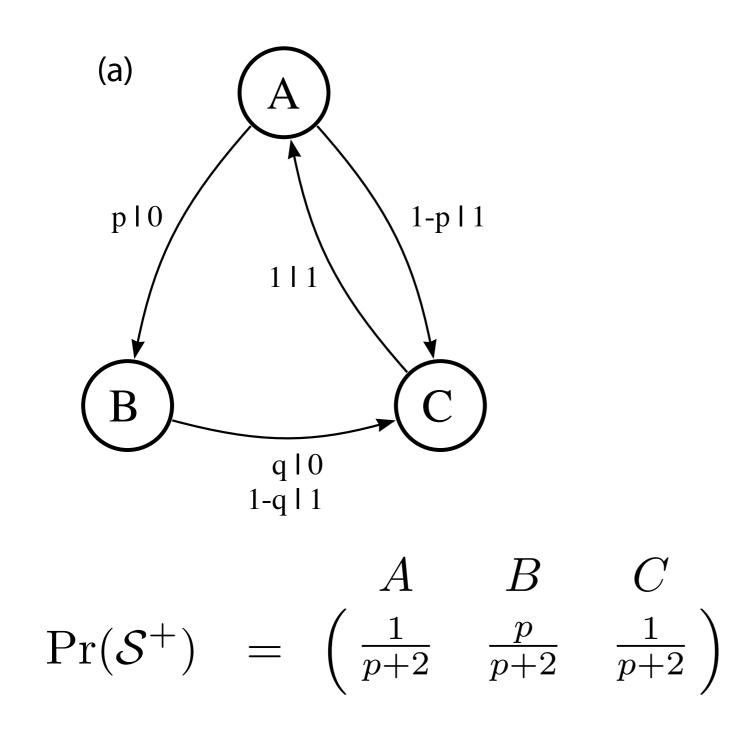
Subshift of finite-type But simple: Causally reversible:  $\Xi = 0$ and not simple: Cryptic process:  $\chi^+ = \chi^- = \frac{H(p)}{2-p}$ 

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Causal Irreversibility

**Example: Random Insertion Process** 

Forward ε-machine:

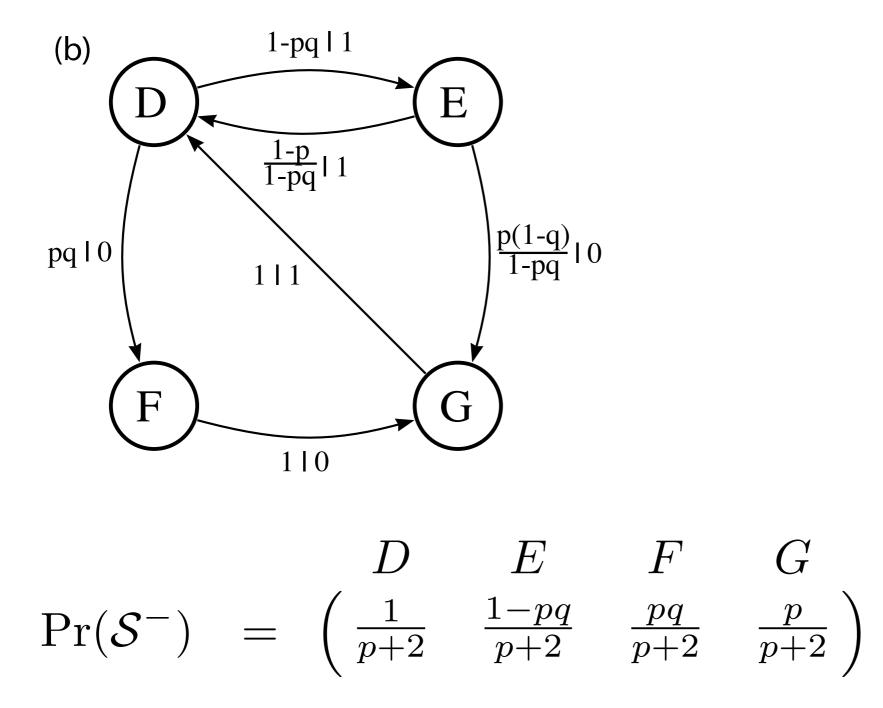


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# Causal Irreversibility

# Example: Random Insertion Process ...

# Reverse E-machine:



# Example: Random Insertion Process ...

# Causally irreversible:

$$M^{+} \neq M^{-}$$

$$C_{\mu}^{+} = \log_{2}(p+2) - \frac{p \log_{2} p}{p+2}$$

$$C_{\mu}^{-} = \log_{2}(p+2) + \frac{H(pq) - p \log_{2} p}{p+2}$$

$$C_{\mu}^{+} \neq C_{\mu}^{-}$$
$$\Xi = \frac{H(pq)}{p+2}$$

Example: Random Insertion Process ...

Forward switching map:  $p = q = \frac{1}{2}$ 

$$\Pr(\mathcal{S}^{+}|\mathcal{S}^{-}) = \frac{D}{F} \begin{pmatrix} 0 & 0 & 1\\ 2/3 & 1/3 & 0\\ 0 & 1 & 0\\ 1 & 0 & 0 \end{pmatrix}$$

Reverse switching map:  $p = q = \frac{1}{2}$ 

$$\Pr(\mathcal{S}^{-}|\mathcal{S}^{+}) = \begin{array}{cccc} B & C & F & G \\ A & \begin{pmatrix} 0 & 1/2 & 0 & 1/2 \\ 0 & 1/2 & 1/2 & 0 \\ 1 & 0 & 0 & 0 \end{array} \right)$$

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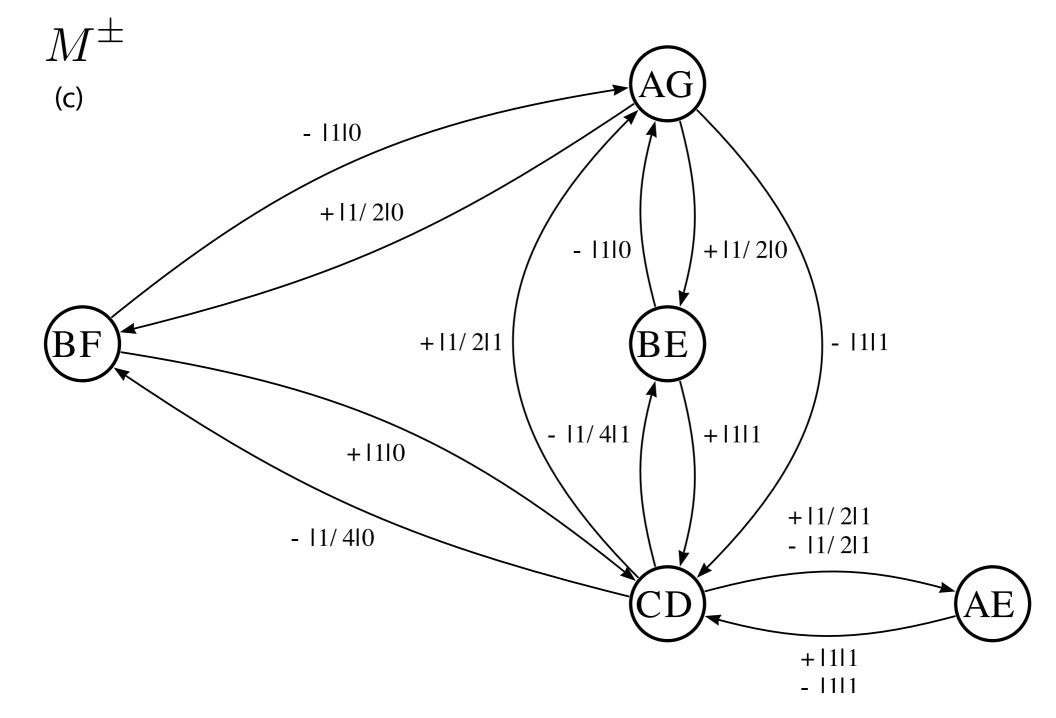
Example: Random Insertion Process ...

Joint distribution (general p and q):

$$\Pr(\mathcal{S}^+, \mathcal{S}^-) = \frac{1}{(p+2)} \frac{A}{C} \begin{pmatrix} 0 & 1-p & 0 & p \\ 0 & p(1-q) & pq & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

Example: Random Insertion Process ...





Example: Random Insertion Process ...

$$\chi^{+} = \frac{1 - pq}{p + 2} H\left(\frac{1 - p}{1 - pq}\right)$$

$$\chi^{-} = \frac{1 - pq}{p + 2} H\left(\frac{1 - p}{1 - pq}\right) + \frac{H\left(pq\right)}{p + 2}$$

$$\mathbf{E} = C_{\mu}^{+} - \chi^{+}$$
  
=  $\log_{2}(p+2) - \frac{p \log_{2} p}{p+2} - \frac{1 - pq}{p+2} H\left(\frac{1 - p}{1 - pq}\right)$ 

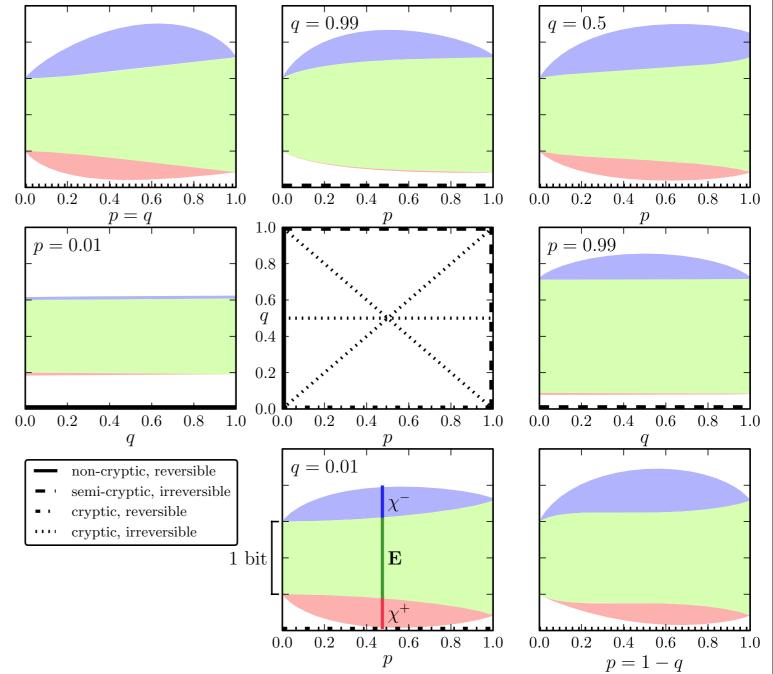
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Example: Random Insertion Process ...

Generally, cryptic, irreversible process  $(p,q) \in [0,1]^2$ 

But ranges over:

non-cryptic, reversible semi-cryptic, irreversible cryptic, reversible cryptic, irreversible



Comments:

Naive combination of forward and reverse machines would be addition or product over states, bidirectional machine is neither.

Bidirectional machine:

- Nonminimal Nonunifilar
- Not E-machine.

Projections onto forward moves gives forward process, but might be nonminimal machine.Projection onto reverse moves gives reverse process, but might be nonminimal machine.

Oddities of "prediction":

Predictive states are better retrodictors than they are predictors (by  $\chi^+$  ).

Predictive states are better retrodictors than retrodictive states (by  $\chi^-$ ).

Open questions:

What are bidirectional transient states?

Minimality?

Unifilar presentation?

Statistical Complexity of Bidirectional Machine:

$$C^{\pm}_{\mu} \equiv H[\mathcal{S}^{\pm}] = H[\mathcal{S}^{+}, \mathcal{S}^{-}]$$

The stored information required to optimally predict *and* retrodict.

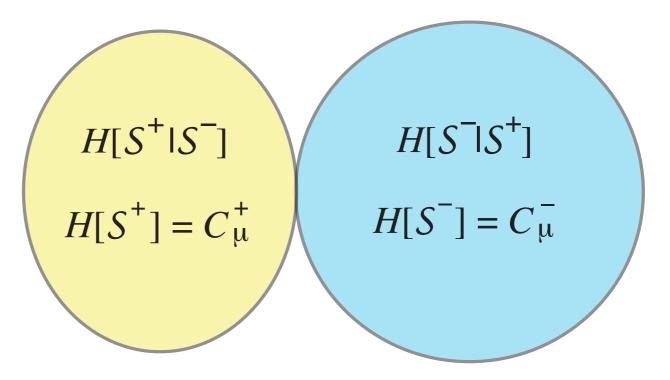
Excess entropy:

$$\mathbf{E} = I[S^{+}; S^{-}]$$
  
=  $H[S^{+}] + H[S^{-}] - H[S^{+}, S^{-}]$   
=  $C_{\mu}^{+} + C_{\mu}^{-} - C_{\mu}^{\pm}$   
 $H[S^{+}] = C_{\mu}^{+}$   
 $H[S^{-}] = C_{\mu}^{+}$ 

Excess entropy ...

Only when  $\mathbf{E} = 0$ :

$$C_{\mu}^{\pm} = C_{\mu}^{+} + C_{\mu}^{-}$$



Bounds:

Before:

$$\mathbf{E} \le C_{\mu} \ (\equiv C_{\mu}^{+})$$

Now, tighter bounds on excess entropy:

$$\mathbf{E} \leq C_{\mu}^+$$
  
and  
 $\mathbf{E} \leq C_{\mu}^-$ 

Bounds ...

Bidirectional machine smaller than forward and reverse:

 $C^{\pm}_{\mu} \leq C^{+}_{\mu} + C^{-}_{\mu}$   $M^{\pm}$  is efficient representation.

Forward E-machine smaller than bidirectional:

 $C^+_{\mu} \le C^{\pm}_{\mu}$ 

Reverse E-machine smaller than bidirectional:

$$C_{\mu}^{-} \le C_{\mu}^{\pm}$$

# From I-diagram:

$$C^{\pm}_{\mu} = \mathbf{E} + \chi^{\pm}$$

#### where

Crypticity: 
$$\chi^{\pm} = \chi^{+} + \chi^{-}$$
  
=  $H[\mathcal{S}^{+}|\mathcal{S}^{-}] + H[\mathcal{S}^{-}|\mathcal{S}^{+}]$ 

Distance between measurements & model:  $C_{\mu}^{\pm} - \mathbf{E}$ 

Form of a true distance: d(X, Y) = H[X|Y] + H[Y|X]

Also, distance between forward and reverse processes.

# Information inaccessibility:

Degree to which internal information is hidden.

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# Bounds ...

$$\chi^{\pm} \le C_{\mu}^{\pm}$$

# Truly cryptic process:

 $\mathbf{E} = \mathbf{0}$ 

$$\chi^{\pm}=C^{\pm}_{\mu}$$
 (All state information is in crypticity.)

# Nothing can be learned about a process's structure from measurements.

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# Summary

Bidirectionality Excess entropy from E-machine Information diagrams for processes Bidirectional machines Bidirectional complexities