

Measures of Structural Complexity

Reading for this lecture:

CMR articles *CMPPSS* and *RURO* and *Lecture Notes*.

Measures of Complexity ...

Information Measures of Complexity:

Measures		Interpretation
Entropy Rate	h_μ	Intrinsic Randomness
Excess Entropy	E	Info: Past to Future
Predictability Gain	G	Redundancy
Transient Information	T	Synchronization

How related to stored information (statistical complexity)?

How to get from ϵM ?

Measures of Complexity ...

Measures from the $\epsilon\mathcal{M}$:

Entropy Rate of a Process:

$$h_{\mu}(\text{Pr}(\overleftrightarrow{\mathcal{S}})) = \lim_{L \rightarrow \infty} \frac{H(L)}{L}$$

Directly from process's $\epsilon\mathcal{M}$:

$$h_{\mu}(\text{Pr}(\overleftrightarrow{\mathcal{S}})) = h_{\mu}(\mathcal{S})$$

Measures of Complexity ...

Measures from the $\epsilon\mathcal{M}$:

Entropy Rate given $\epsilon\mathcal{M}$:

$$h_{\mu}(\mathcal{S}) = - \sum_{\mathcal{S} \in \mathcal{S}} \text{Pr}(\mathcal{S}) \sum_{s \in \mathcal{A}, \mathcal{S}' \in \mathcal{S}} T_{\mathcal{S}\mathcal{S}'}^{(s)} \log_2 T_{\mathcal{S}\mathcal{S}'}^{(s)}$$

where $\text{Pr}(\mathcal{S})$ is casual-state asymptotic probability.

Possible only due to $\epsilon\mathcal{M}$ unifilarity!

I-I mapping between measurement sequences & internal paths.

Measures of Complexity ...

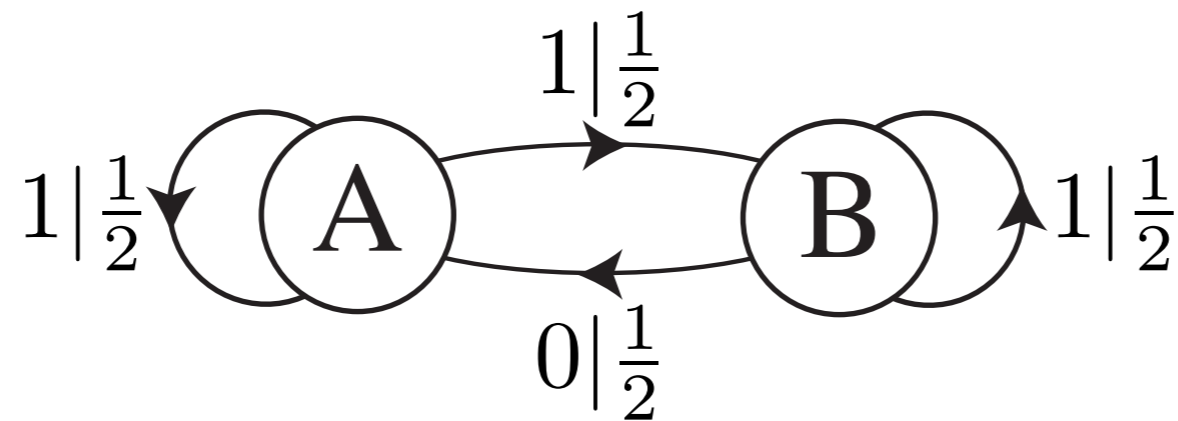
Measures from the ϵM :

Entropy rate ...

Possible only due to unifilarity ... Why?

What if you use a nonunifilar HMM presentation?

Consider nonunifilar presentation of SNS Process:



$$\mathcal{B} = \{0, 1\}$$

$$T^{(0)} = \begin{pmatrix} 0 & 0 \\ \frac{1}{2} & 0 \end{pmatrix} \quad T^{(1)} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ 0 & \frac{1}{2} \end{pmatrix}$$

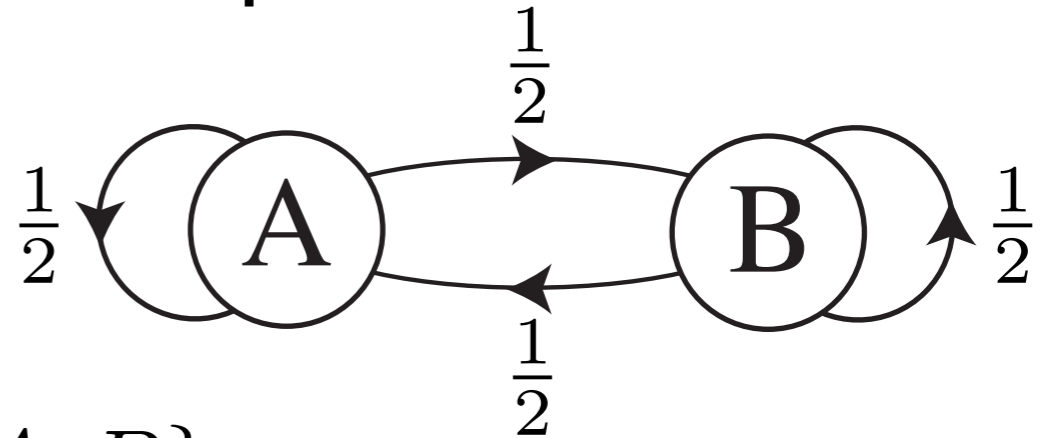
Measures of Complexity ...

Measures from the ϵM :

Method A:

Calculate entropy rate from nonunifilar presentation:

Internal Markov chain:



Internal (= Fair Coin): $V = \{A, B\}$

$$T = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{pmatrix} \quad \pi_V = \left(\frac{1}{2}, \frac{1}{2} \right)$$

Entropy rate?

$$\begin{aligned} \text{“}h_{\mu}^A(\text{SNS})\text{”} &= - \sum_{v \in V} \Pr(v) \sum_{v' \in V} \Pr(v'|v) \log_2 \Pr(v'|v) \\ &= 1 \text{ bit/symbol} \end{aligned}$$

Measures of Complexity ...

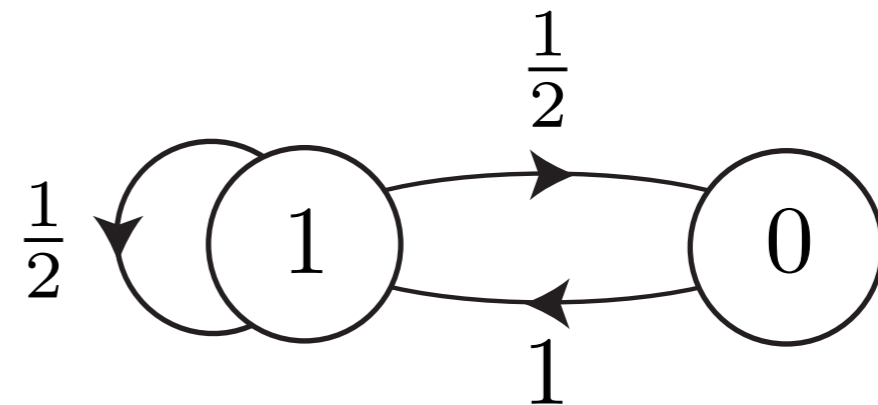
Measures from the $\epsilon\mathcal{M}$:

Method B:

Lower bound by entropy rate for “support” process:

No consecutive 0s!

Support process is GMP



$$T = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ 1 & 0 \end{pmatrix} \quad \pi_V = (2/3, 1/3)$$

A restriction that lowers entropy rate:

$$\begin{aligned} h_{\mu}^B(\text{GMP}) &= -\frac{2}{3} \sum_{s \in \mathcal{B}} \text{Pr}(s|1) \log_2 \text{Pr}(s|1) + 0 \\ &= \frac{2}{3} \approx 0.666667 \text{ bits/symbol} \end{aligned}$$

Measures of Complexity ...

Measures from the $\epsilon\mathcal{M}$:

“SNS entropy rate” larger than support process?

$$“h_{\mu}^A(\text{SNS})” \gg h_{\mu}^B(\text{GMP})$$

Not possible!

Entropy of distribution always \leq Entropy of support

Use of SNS presentation leads to overestimate of entropy rate:
internal process more random than observed process.

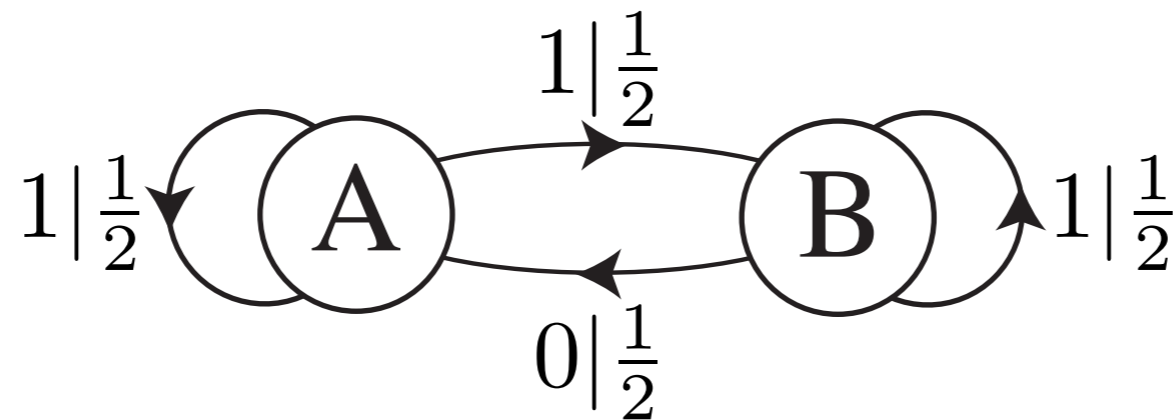
So how to compute SNS process’s entropy rate?

Measures of Complexity ...

Measures from the $\epsilon\mathcal{M}$:

Method C:

Estimate entropy rate from output process:



$$\begin{aligned} \text{“}h_{\mu}^C(\text{SNS})\text{”} &= - \sum_{v \in V} \Pr(v) \sum_{(s, v') \in (\mathcal{A}, V)} \Pr(s, v' | v) \log_2 \Pr(s, v' | v) \\ &= - \left(\frac{1}{2} \times 0 - \frac{1}{2} \times 1 \right) \\ &= \frac{1}{2} \text{ bits/symbol} \end{aligned}$$

Measures of Complexity ...

Measures from the ϵM :

Entropy rate?

$$“h_{\mu}^A(\text{SNS})” \gg h_{\mu}^B(\text{GMP}) \gg “h_{\mu}^C(\text{SNS})”$$

Lessons:

1. Cannot use nonunifilar HMM presentations for entropy rate.
2. Need ϵM to calculate entropy rate.

SNS example: Nontrivial, countably infinite ϵM .

$$h_{\mu}(\text{SNS}) \approx 0.6778 \text{ bits/symbol}$$

Curious:

Even to estimate a process's intrinsic randomness,
need to infer its structure.

Measures of Complexity ...

Measures from the ϵM ...

Statistical Complexity of ϵM :

$$C_{\mu}(\mathcal{S}) = - \sum_{\mathcal{S} \in \mathcal{S}} \text{Pr}(\mathcal{S}) \log_2 \text{Pr}(\mathcal{S})$$

where $\text{Pr}(\mathcal{S})$ is casual-state asymptotic probability.

Meaning:

Shannon information in the causal states.

Measures of Complexity ...

Measures from the ϵ M...

Statistical Complexity of a Process:

$$C_{\mu} \left(\text{Pr}(\vec{S}) \right) = C_{\mu}(\mathcal{S})$$

Meaning:

The amount of historical information a process stores.

The amount of structure in a process.

Measures of Complexity ...

Measures from the $\epsilon\mathcal{M}$...

Excess Entropy: Three versions, all equivalent for IID processes

$$\mathbf{E} = \lim_{L \rightarrow \infty} [H(L) - h_\mu L]$$

$$\mathbf{E} = \sum_{L=1}^{\infty} [h_\mu(L) - h_\mu]$$

$$\mathbf{E} = I[\overleftarrow{S}; \overrightarrow{S}]$$

How to get, given $\epsilon\mathcal{M}$?

Special cases: When $\epsilon\mathcal{M}$ is IID, periodic, or spin chain.

General case: Need a new framework ... next lectures.

Measures of Complexity ...

Measures from the ϵM ...

Excess Entropy ...

ϵM is IID:

$$\mathbf{E} = 0 \quad C_\mu = 0$$

ϵM is Period P :

$$\mathbf{E} = \log_2 P = C_\mu$$

ϵM is range- R spin chain:

$$\mathbf{E} = H(R) - Rh_\mu$$

$$\mathbf{E} = C_\mu - Rh_\mu$$

Measures of Complexity ...

Measures from the ϵ M...

Excess Entropy ...

Typically for Markov Chains:

$$\mathbf{E} < C_{\mu}$$

What can be said in general?

Measures of Complexity ...

Measures from the $\epsilon\mathbb{M}$...

Bound on Excess Entropy:

$$\mathbf{E} \leq C_\mu$$

Proof sketch:

$$(1) \mathbf{E} = I[\vec{\mathcal{S}}; \overleftarrow{\mathcal{S}}] = H[\vec{\mathcal{S}}] - H[\vec{\mathcal{S}} | \overleftarrow{\mathcal{S}}]$$

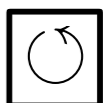
$$(2) \text{ Causal States: } H[\vec{\mathcal{S}} | \overleftarrow{\mathcal{S}}] = H[\vec{\mathcal{S}} | \mathcal{S}]$$

$$(3) \mathbf{E} = H[\vec{\mathcal{S}}] - H[\vec{\mathcal{S}} | \mathcal{S}]$$

$$= I[\vec{\mathcal{S}}; \mathcal{S}]$$

$$= H[\mathcal{S}] - H[\mathcal{S} | \vec{\mathcal{S}}]$$

$$\leq H[\mathcal{S}] = C_\mu$$



Measures of Complexity ...

Measures from the $\epsilon\mathbb{M}$...

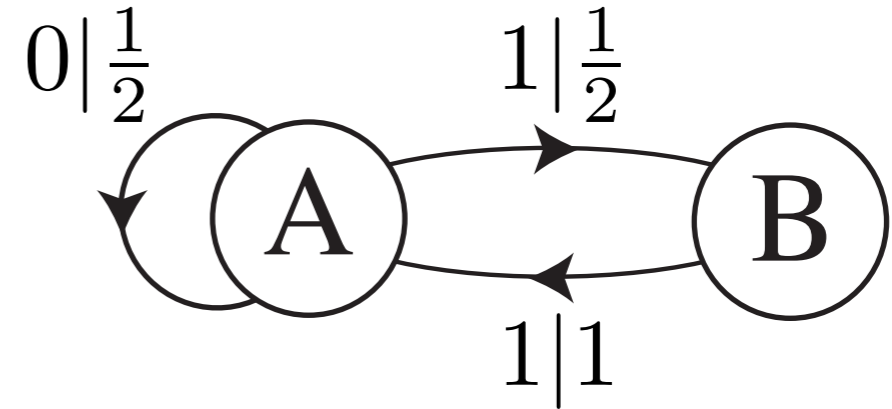
Bound on Excess Entropy ...

But, the bound is saturated!

Even process:

$$C_\mu = H(2/3) \approx 0.9182$$

$$\mathbf{E} \approx 0.9182$$



$$T^{(0)} = \begin{pmatrix} \frac{1}{2} & 0 \\ 0 & 0 \end{pmatrix}$$

$$T^{(1)} = \begin{pmatrix} 0 & \frac{1}{2} \\ 1 & 0 \end{pmatrix}$$

$$\pi_V = (2/3, 1/3)$$

When does this occur?

In general, need a new framework for answering this question.

Measures of Complexity ...

Measures from the ϵM ...

Bound on Excess Entropy ...

Consequence: **The Cryptographic Limit**

Can have $\mathbf{E} \rightarrow 0$ when $C_\mu \gg 1$.

State information is still there,
just hidden in the observed sequences.

Goal of cryptography:

Cyphertext appears random to all observers,
but still contains the original information,
recoverable by the intended recipient.

Measures of Complexity ...

Measures from the ϵ M ...

Bound on Excess Entropy ...

For example: The Almost IID Processes

$$p_i \sim \frac{1}{2} + \text{Uniform}(0, \epsilon)$$

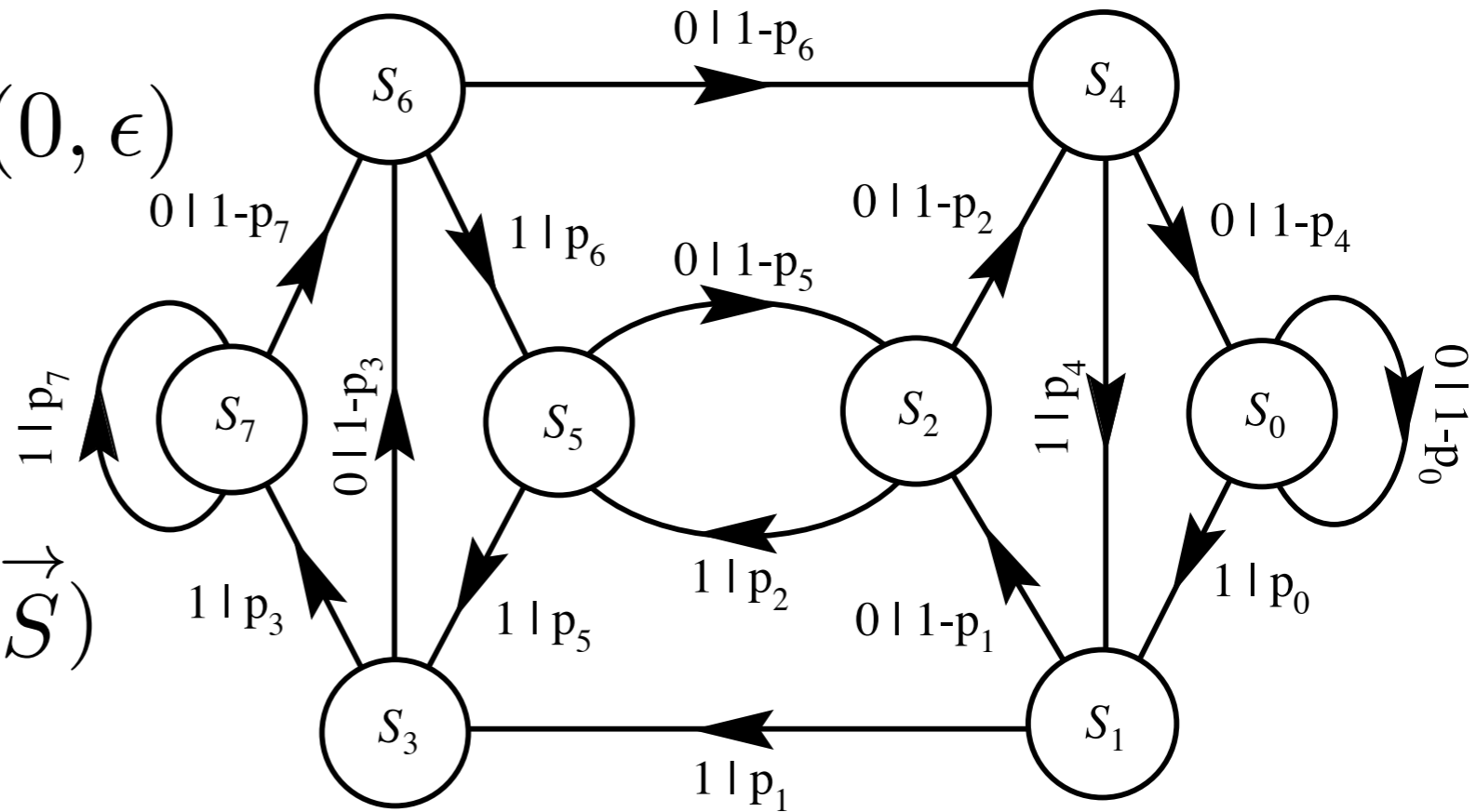
$\epsilon \searrow 0$

$$\Pr(\vec{S}) \approx \Pr(\overleftarrow{S})\Pr(\overrightarrow{S})$$

$$\mathbf{E} \approx 0$$

$$\Pr(\mathcal{S}_i) \approx |\mathcal{S}|^{-1}$$

$$C_\mu \approx \log_2 |\mathcal{S}|$$



Measures of Complexity ...

Measures from the ϵ M ...

Bound on Excess Entropy ... $\mathbf{E} \leq C_\mu$

Consequence:

Excess entropy is *not* the process' stored information.

Statistical complexity C_μ is the stored information.

\mathbf{E} is the *apparent* information,
as revealed in *measurement sequences*.

Measures of Complexity ...

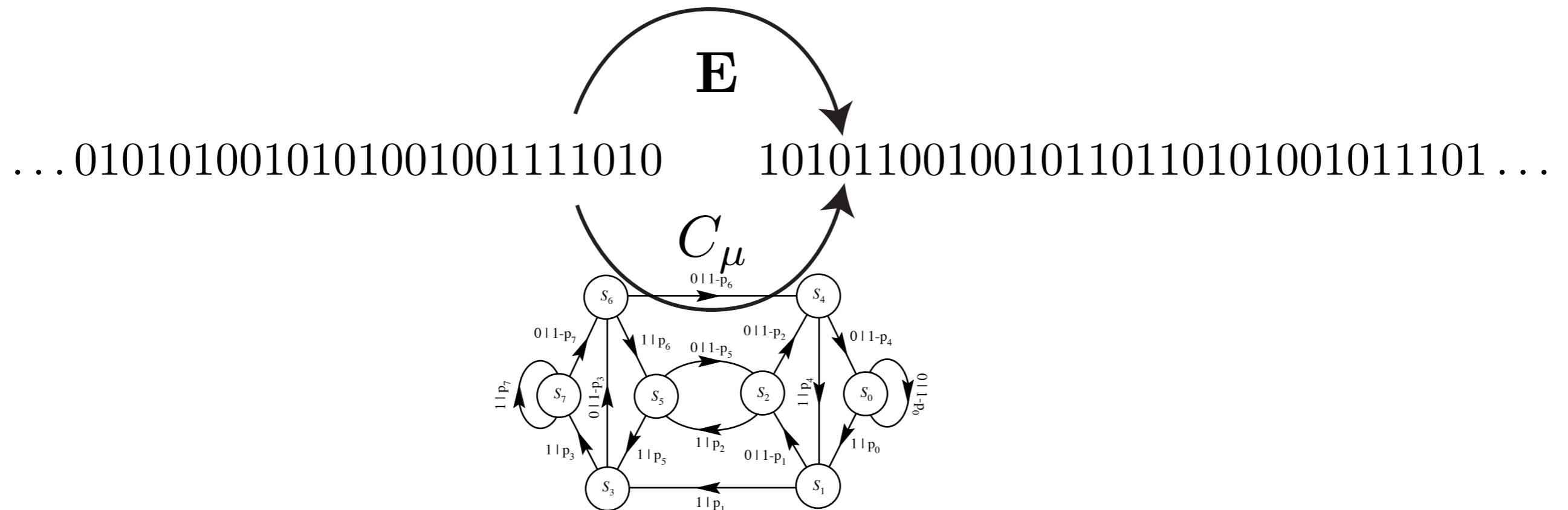
Measures from the ϵM ...

Bound on Excess Entropy ...

Executive Summary:

C_μ is the amount of information the process uses
to *communicate*

E bits of information from the past to the future.



Measures of Complexity ...

Measures from the $\epsilon\mathbb{M}$...

Bound on Excess Entropy: $\mathbf{E} \leq C_\mu$

Consequence:

The inequality is Why We Must Model.

Cannot simply use sequences as states.

There is internal structure not expressed by sequences.

Widespread misconception about excess entropy $I[\overleftarrow{X}; \overrightarrow{X}]$:
“the predictive information”.

No! C_μ is the information required to optimally predict.

Excess entropy is the “predictable” information:
that which can be predicted.

Measures of Complexity ...

Measures from the $\epsilon\mathbb{M}$...

Bound on Predictability (Redundancy):

$$r(1) \leq C_\mu$$

Proof sketch:

$$(1) \quad r(1) = H[\overrightarrow{S}^1] - h_\mu$$

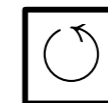
$$(2) \quad H[\overrightarrow{S}^1] - h_\mu = H[\overrightarrow{S}^1] - H[\overrightarrow{S}^1 | \overleftarrow{S}]$$

$$= H[\overrightarrow{S}^1] - H[\overrightarrow{S}^1 | \mathcal{S}]$$

$$= I[\overrightarrow{S}^1; \mathcal{S}]$$

$$\leq H[\mathcal{S}]$$

$$= C_\mu$$



Measures of Complexity ...

An ϵM Captures “Pattern”:

Symmetry = group:

Rotate square 90° in plane.

Square appears in original orientation.

Square “shape” is invariant under 90° rotation.

Inverse: Rotate by -90° .

Measures of Complexity ...

An $\epsilon\mathbb{M}$ Captures “Pattern” ...

$\epsilon\mathbb{M}$ is a **monoid**:

$\{T_{ij}^{(s)}\}$ form a semi-group with identity.

Semi-group is a group without a unique inverse.

When semi-group is a group: $\epsilon\mathbb{M}$ algebra describes a symmetry.

Semi-groups describe “generalized”, noisy, wild-card symmetries.

Measures of Complexity ...

An ϵM Captures “Pattern”:

Example of Group: Period-2 process

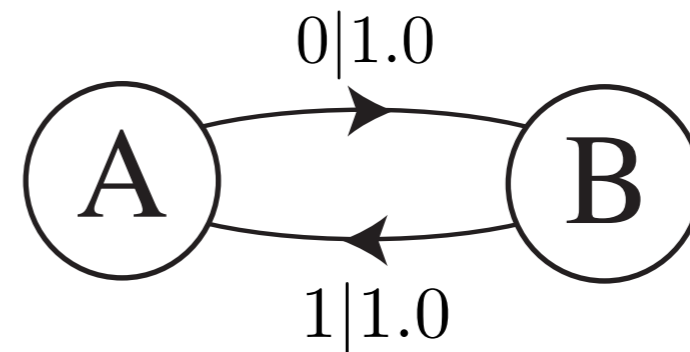
... 01010101010101010101 ...

... 01010101010101010101 ...

Temporal translation symmetry:

Shift time by two steps ($t \rightarrow t + 2$), get same sequence.

ϵM period-2 in the causal states:



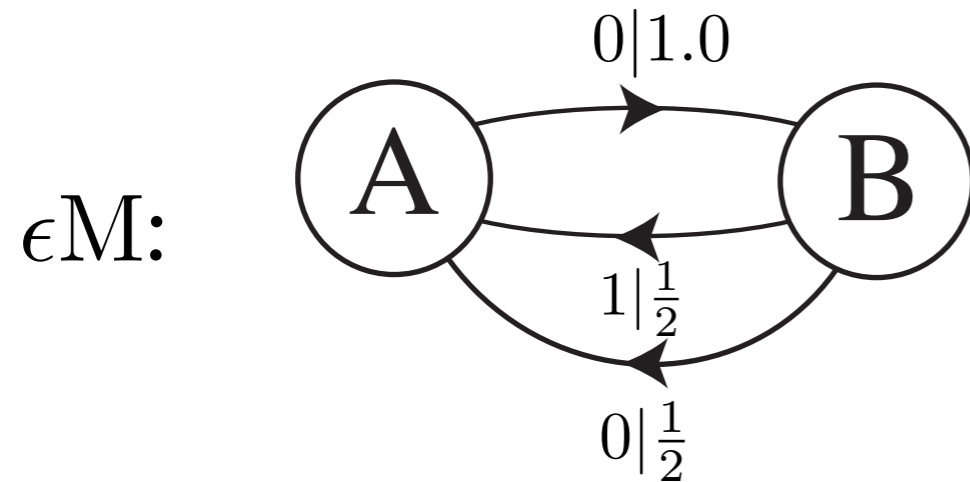
... *ABABABABABAB* ...

Measures of Complexity ...

An ϵM Captures “Pattern”:

Example: Random process

... 010100010001010100010001 ...



Period-2 in the causal states: Detects the order in chaos

... *ABABABABABAB* ...

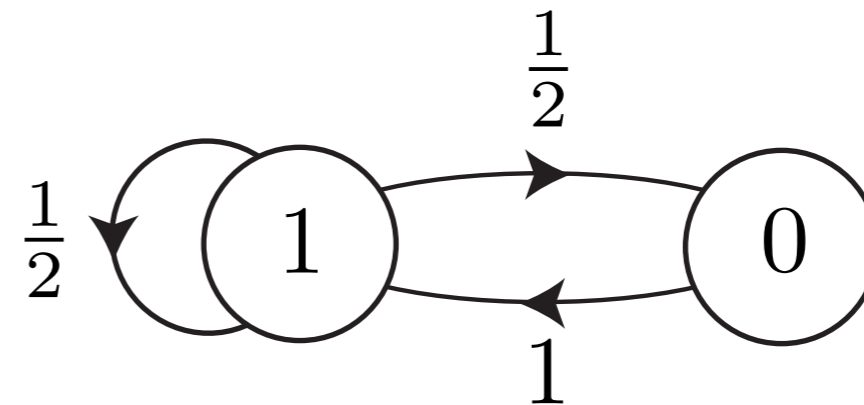
Temporal translation symmetry:

Shift time by two steps ($t \rightarrow t + 2$), get **same statistics**.

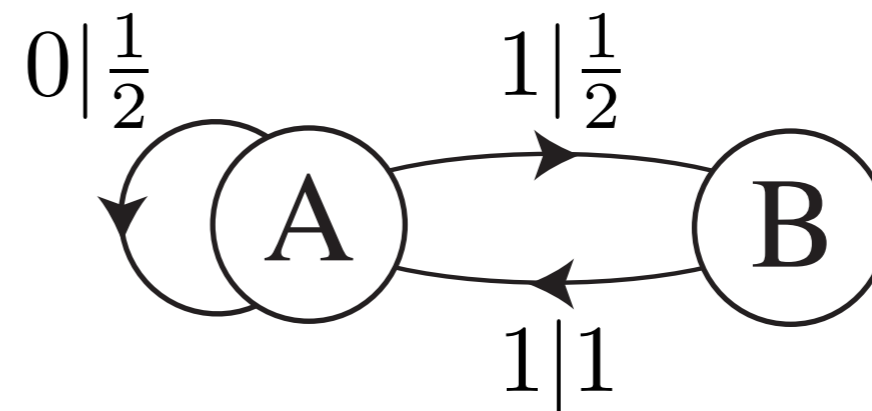
Measures of Complexity ...

An ϵM Captures “Pattern”:

Example: GMS “pattern” = No consecutive 0s
By inspecting ϵM structure.



Example: Even process “pattern” = “Evenness”
By inspecting ϵM structure.

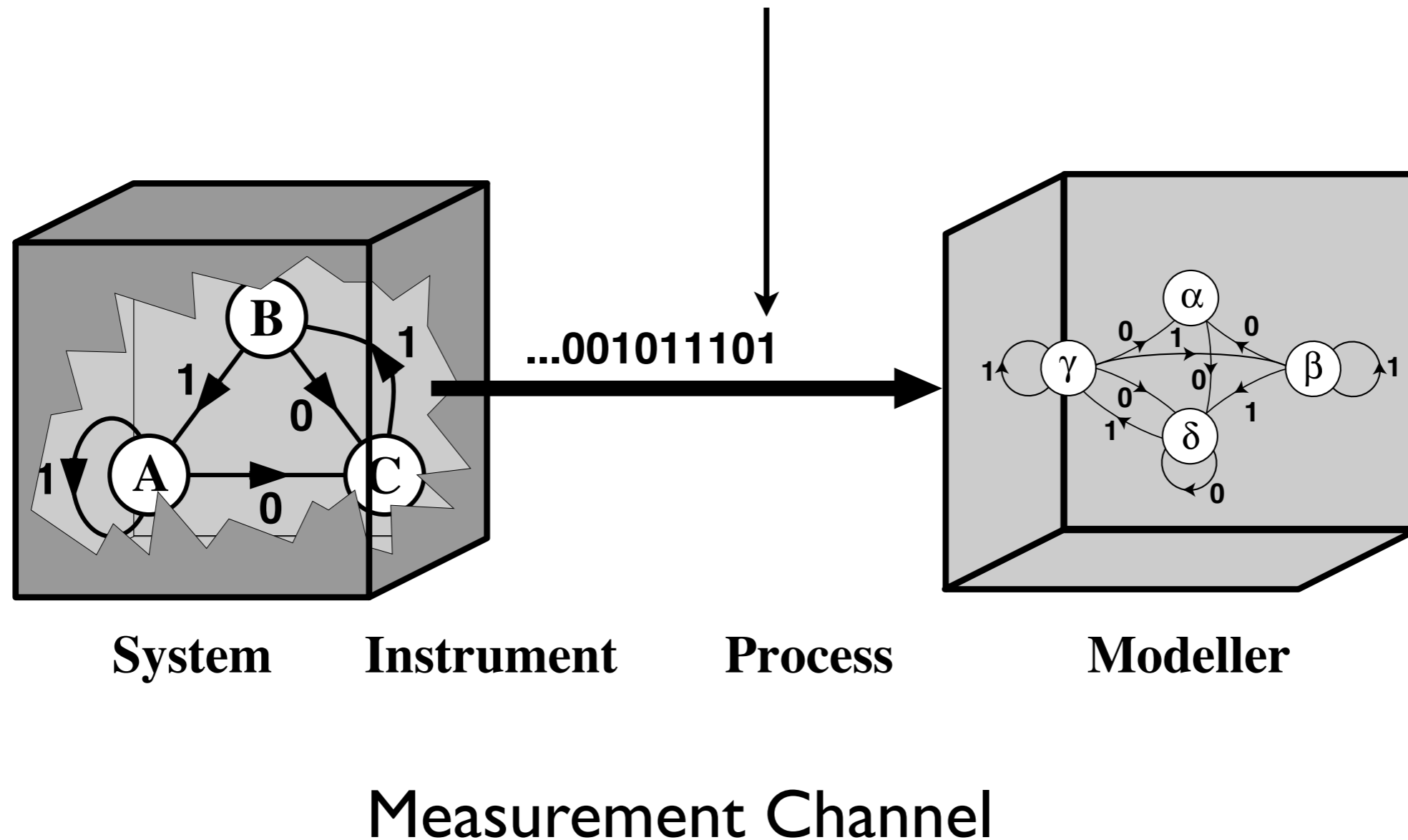


Measures of Complexity ...

An ϵM Captures “Pattern”:

Measurement semantics:

What does a particular measurement mean?



Measures of Complexity ...

An ϵ M Captures “Pattern”:

Measurement semantics: Prediction level

What is the meaning of a particular measurement s ?

Shannon says the amount of “information” in s is:

$$-\log_2 \Pr(\text{observing } s)$$

Given ϵ M (assuming you’re sync’d to internal (causal) state):

$$-\log_2 \Pr(\text{observing } s) = -\log_2 \Pr(\mathcal{S} \rightarrow_s \mathcal{S}')$$

Measures of Complexity ...

An ϵ M Captures “Pattern”:

Measurement semantics: Prediction level ...

Example: $t = 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11$

$s = 0 \ 1 \ 1 \ 1 \ 1 \ 0 \ 1 \ 1 \ 0 \ 1 \ 1 \ 1$

At $t = 11$ measure $s_{11} = 1$

How much information does this give?

$$H(s_{11}|s_{10} = 1, s_9 = 1, \dots) \approx h_\mu(\approx 0.585 \text{ bits})$$

Degree of observer’s surprise (predictability)

Does not say what the event $s_{11} = 1$ means to the observer!

Measures of Complexity ...

An ϵM Captures “Pattern”:

Measurement semantics ...

Meaning: Tension between representations of same event at different levels; e.g.:

Level 1 is data stream and the event is a measurement

Level 2 is the agent and the event updates it's model

Degree of meaning of observing $s \in \mathcal{A}$:

$$\Theta(s) = -\log_2 \Pr(\rightarrow_s \mathcal{S})$$

where \mathcal{S} is the causal state to which s brings observer.

Meaning content: State selected from anticipated palette.

Measures of Complexity ...

An ϵ M Captures “Pattern”:

Measurement semantics ...

Meaningless: Start state (all futures possible)

$$\Theta(s) = -\log_2 \Pr(\mathcal{S}_0) = -\log_2 1 = 0 \quad s = \lambda$$

Action on disallowed transition:

Reset to state of total ignorance (start state)

Disallowed transition is meaningless.

Meaningless measurements are informative, though:

$$-\log_2 \Pr(\mathcal{S} \rightarrow_s \mathcal{S}_0) = -\log_2 0 = \infty$$

Measures of Complexity ...

An ϵ M Captures “Pattern”:

Measurement semantics ...

Theorem:

$$\langle \Theta(s) \rangle = C_\mu$$

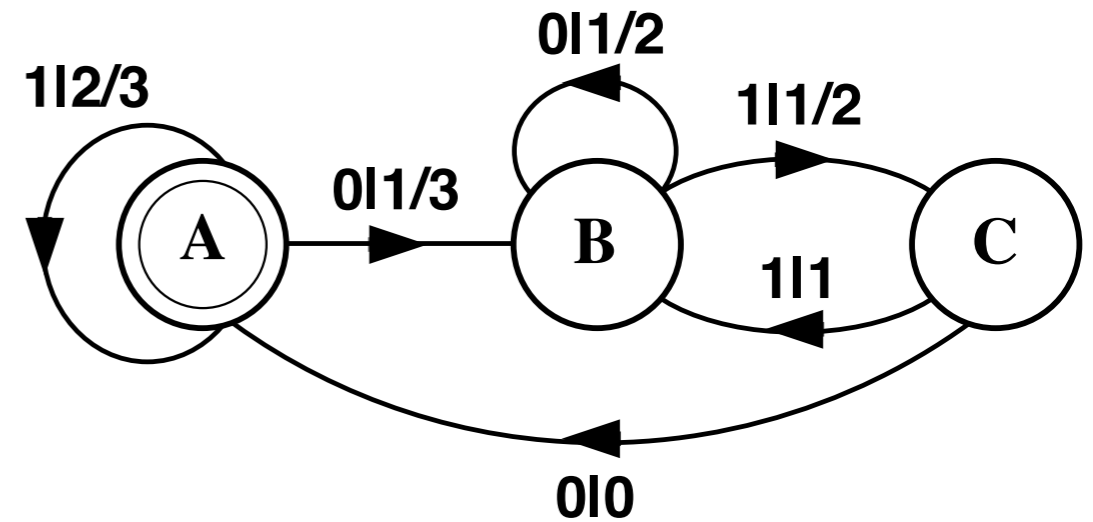
Average amount of meaning is the Statistical Complexity.

Measures of Complexity ...

An ϵM Captures “Pattern”:

Measurement semantics: Example

...110110111101111101101...



Observer's Semantic Analysis				
State	Measurement	Surprise (bits)	Semantic State: Meaning	Degree of Meaning [bits]
A	λ	Not defined	No measurement	0
A	1	0.585	A: Unsynchronized	0.585 ... ∞
A	0	1.585	B: Synchronized	1.585 ... 0.585
B	1	1	C: Odd # 1s	1.585
B	0	1	B: Even # 1s	0.585
C	1	0	B: Even # 1s	0.585
C	0	∞	A: Loose sync; reset	0

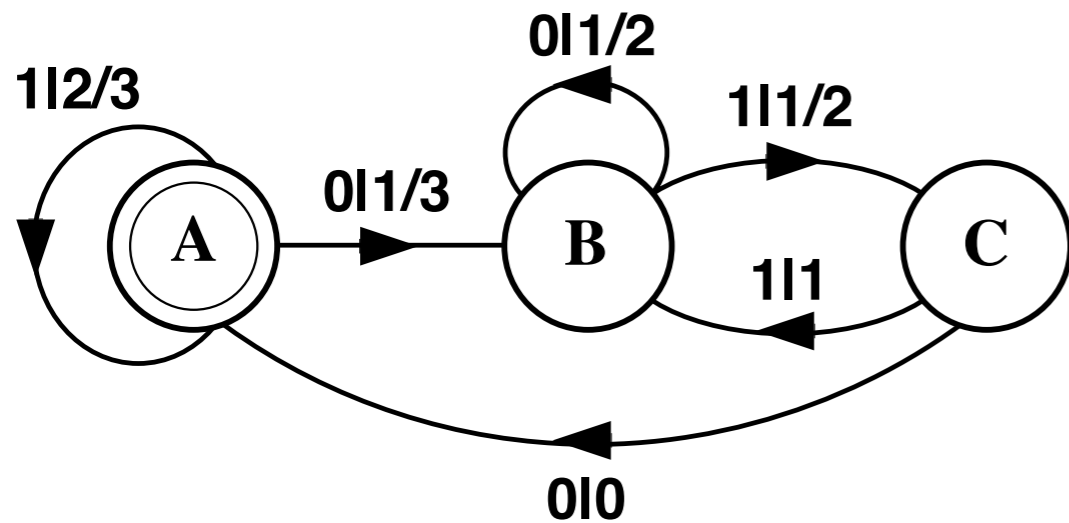
... the Even Process

Measures of Complexity ...

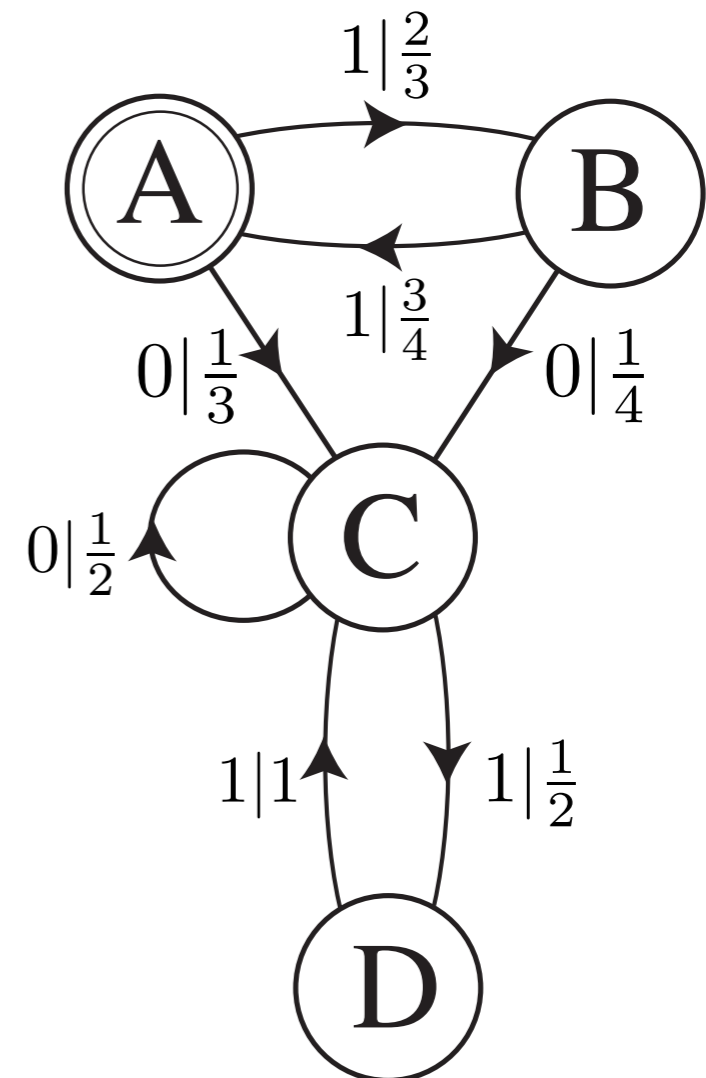
An ϵ M Captures “Pattern”:

Measurement semantics:

... the Even Process?



No, not exactly:



Semantics induced by any model:
Even an incorrect one!

Measures of Complexity ...

An ϵM Captures “Pattern”:

Measurement semantics: Slipped something by ...

Preceding was about the *amount* of meaning!

Related to amount of memory/structure

But what is the meaning content?

Meaning = Algebra of the ϵM

Measures of Complexity ...

An ϵM Captures “Pattern”:

Measurement semantics ...

D. M. MacKay, “Information, Meaning & Mechanism”, MIT (1969):

Meaning as selection of anticipated context.

Here:

ϵM s give the natural set of contexts (causal states)
for an information source.

Measures of Complexity ...

Reading for next lecture:

CMR articles TBA and PRATISP.